

# Resilience in Information Stewardship

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**Abstract.** Information security is concerned with protecting the confidentiality, integrity, and availability of information systems. System managers deploy their resources with the aim of maintaining target levels of these attributes in the presence of reactive threats. Information stewardship is the challenge of maintaining the sustainability and resilience of the security attributes of (complex, interconnected, multi-agent) information ecosystems. In this paper, we present, in the tradition public economics, a model of stewardship which addresses directly the question of resilience. We model attacker-target-steward behaviour in a fully endogenous Nash equilibrium setting. We analyse the occurrence of externalities across targets and assess the steward's ability to internalize these externalities under varying informational assumptions. We apply and simulate this model in the case of a critical national infrastructure example.

## 1 Introduction

The objective of information security managers is to protect the confidentiality, integrity, and availability of the information systems for which they are responsible, by adjusting their adopted security measures in response to the evolving threat environment. The optimal investment in such security measures has been studied initially by Gordon and Loeb [8].

The dynamic responses of information security managers to evolving threats to confidentiality, integrity, and availability have been modelled by, among others, Ioannidis, Pym, and Williams [10], August, and Tunca [3] and Arora, Telang and Xu [1].

This objective must be pursued in the context of an information ecosystem that is subject to finite degradation of performance because of internal and external influences. In the information ecosystem, threats to the confidentiality, integrity, and availability of individual components of the ecosystem can be transmitted to others, impacting negatively on their security status. In such an environment, the role of the information steward is to maintain the sustainability

and resilience of the ecosystem's nominal operating capacity, so delivering the managers' desired levels of confidentiality, integrity, and availability.

### 1.1 Sustainability

By the sustainability of a system, subject to finite degradation caused by a persistent stream of attacks, we mean its tendency to remain within specified levels of nominal operating capacity. The graph in Figure 1 depicts this concept in static framework.

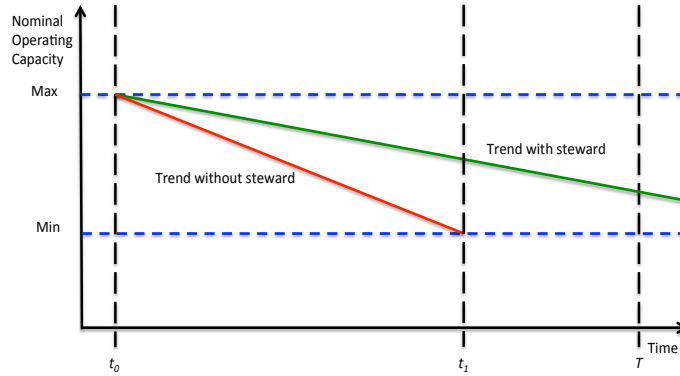


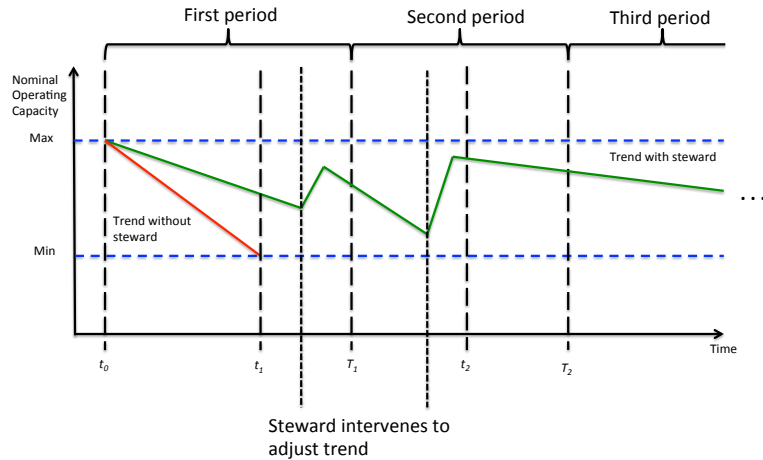
Fig. 1. The concept of sustainability in information stewardship.

During successive tenures, stewards adopt policies to extend the system's lifetime. This is illustrated in Figure 2, in which interventions by the steward can be seen to maintain the system within its intended operating zone. From a game-theoretic point of view, managing the ecosystem to achieve this objective is a mechanism design problem for the steward.

In a recent paper, Ioannidis, Pym, and Williams [10] have explored the steward's responses to maintain sustainability in the presence of endogenous attacks. They show that the presence of the steward increases investment in information security and, more importantly, reduces the number of attacks, thus retarding system degradation.

### 1.2 Resilience

In addition to the persistent stream of attacks to which a system is regularly subjected, occasionally the impact of an attack will be such that the system diverges from its predefined operational bounds. By resilience, we mean the ability of the system to return its operating capacity to within the specified bounds. Our notion of resilience is consistent with that discussed in [9, 15].



**Fig. 2.** A multi-stage sustainability problem with periodic technology resets.

The steward seeks to minimize, subject to a range of endogenous constraints, the time for which the system operates outside of the specified operating limits.

Figure 3 depicts the concept of resilience in static framework in the presence of the information steward.

The remainder of this paper is organized as follows: in § 2, we present a model of resilience, with and without the information steward, with varying degrees of influence over the inhabitants of the ecosystem; in § 3, we present a detailed example of the model, illustrated by simulations, and informed by the operations of Industrial Control Systems (ICSs); finally, § 4 provides a summary of our contribution and some directions for future research.

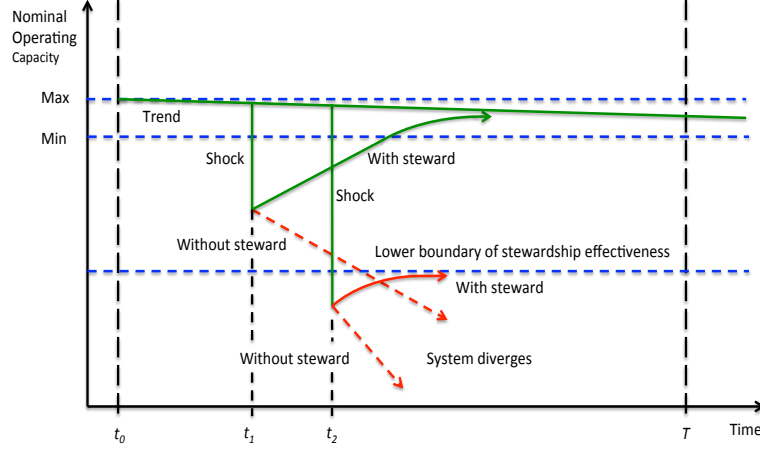
## 2 The Model

We have argued that the challenge for the information steward is to maintain the sustainability and resilience of target levels of confidentiality, availability, integrity, and investment for information ecosystems operating in potentially hostile environments. In previous work [13] and [10], we have shown how to model sustainability, placing our discussion in the broader context of stewardship. We now extend our treatment of sustainability to account for resilience.

This extension is not merely incremental. Rather, it introduces significant additional and conceptual challenges.

A key feature of our account of resilience is that we illustrate how thresholds for the effectiveness of stewards emerge from the underlying model of the response of information ecosystems to the hostility of the environment.<sup>4</sup>

<sup>4</sup> Note by that the hostility of the environment we mean a representation of the capacity of attackers rather than simply the success or failure of an individual attack.



**Fig. 3.** The concept of resilience in information stewardship.

As in our previous work on sustainability [10], our approach is one of mechanism design in which the targets of attacks are expected-loss-minimizers, subject to diminishing marginal returns on security investment. Similarly, attackers are modelled as rational agents. They are assumed to have utility functions, with well defined preferences, which can be used to capture their behavioural choices with respect to a variety of consumption goods, which maybe converted to monetary certainty equivalents, gained from successful attacks. We consider a set of  $N_T$  ex-ante identical targets choosing to allocate defensive resources that mitigate the harm from attacks. In a departure from previous models, the targets need to solve, simultaneously, a multi-dimensional resource allocation problem. Let the subscripts  $h$  and  $l$  represent to potential areas of allocation of assets, where  $h$  and  $l$  denote the areas of high and low security where information assets are held, and let  $x_h \geq 0$  and  $x_l \geq 0$  denote the one-off investments made at time  $t_0$  in securing assets located in the corresponding areas. Finally, we define  $z$  to be a switching variable such that a fraction,  $0 \leq z \leq 1$ , of assets is allocated between  $h$  and  $l$ .

Our model depends crucially on two key (vector) parameters. First, we consider the elasticity of attacking intensity denoted by the vector  $\alpha$ . This is the parameter that captures the marginal effectiveness of an additional attacker per target ( $\eta$ ) entering the ecosystem. Whilst, in general, one would consider an ecosystem of participants having  $n$  types of assets of interest, without loss of representation of resilience, our model restricts is limited to two types of assets. In this case, we need to consider just two elasticities,  $\alpha_l$  and  $\alpha_h$  with corresponding  $\eta_l$  and  $\eta_h$ , which are associated with low and high levels of difficulty in securing assets. Second, we consider parameters  $\psi_l$  and  $\psi_h$ , which capture the relative rate of risk reduction for additional security investments by targets in each asset class ( $x_l$  and  $x_h$ ).

Let  $\tilde{\sigma}_{i \in \{l, h\}} : \mathbb{R}_+ \rightarrow [0, 1]$  be a function that determines the instantaneous time  $t$  risk for a fixed time-horizon, where  $(t_0, T) = \{t \mid t_0 < t < T\}$ . When properly specified we can interpret  $\tilde{\sigma}$  as the instantaneous probability of a successful attack. We refer to  $z$  as the ‘asset allocation’ and the quantities  $x_l$  and  $x_h$  as the ‘investment allocation’, stated combinations of all three are referred to as ‘allocation bundles’.

Our assumption is that increased investment  $x_{i \in \{l, h\}}$  reduces the probability of a successful attack; that is,  $\partial \tilde{\sigma}_{i \in \{l, h\}} / \partial x_{i \in \{l, h\}} < 0$ , ceteris paribus. However, along with increasing investment there is a decreasing marginal reduction in the probability of a successful attack,  $\partial^2 \tilde{\sigma}_{i \in \{l, h\}} / \partial x_{i \in \{l, h\}}^2 > 0$ . Similarly, with increased attacking intensity  $\eta_{i \in \{l, h\}}$  on the particular area of allocation there should be a corresponding increase in the probability of a successful attack,  $\partial \tilde{\sigma}_{i \in \{l, h\}} / \partial \eta_{i \in \{l, h\}} > 0$ .

A functional form for  $\tilde{\sigma}$  that satisfies these conditions is the following multiplicative model:

$$\tilde{\sigma}_i = e^{-\psi_i x_i \eta_i^{\alpha_i}}, \quad i \in \{l, h\}. \quad (1)$$

Under this formulation, there is an upper bound on  $\eta_{i \in \{l, h\}}$  of  $\eta_i^* < e^{\alpha_i^{-1} x_i \psi_i}$ , for  $i \in \{l, h\}$ , such that  $\tilde{\sigma}_i$  may still be interpreted as probability of a successful attack. Here  $\psi_{i \in \{l, h\}}$  is the relative marginal decrease in  $\tilde{\sigma}_{i \in \{l, h\}}$  for a unit increase in  $x_{i \in \{l, h\}}$  whilst  $\alpha_{i \in \{l, h\}}$  is the elasticity of attack.

In this model, we assume that attacker externalities are driven by the diffuse-attacking-mass approach first suggested in [13] and refined in [10]. In this approach attackers are assumed to be ex-ante identical and randomly allocated to targets with identical probability  $1/N_T$ . Attackers are assumed to be able to make independent decisions on the type of attacks. A useful interpretation of the attacker cost per unit is that attackers need to develop an attacking tool at cost  $c$  each time they engage a target. The attacker then chooses the medium by which it seeks to monetize its attacking effort (in the case of terrorists, for example, monetization is via utility equivalents). An example could be corporate network information channels versus industrial control systems. Attackers, at inception, may not know which target they intend to attack. From the viewpoint of the steward in this setting, it is irrelevant who is attacking the targets. From the target-attacker transaction point of view, the salient point is the aggregate level of loss incurred in the presence of attacking intent.

Let the number of attackers for each asset area be  $N_{A, i \in \{l, h\}}$ . The ratio of attackers per target is the attacking intensity  $\eta_{i \in \{l, h\}} = N_{A, i \in \{l, h\}} / N_T$ . Let the reward  $R > 0$  for a successful attack be proportional to the assets allocated in each area,  $h$  and  $l$ , and for notational simplicity let  $\zeta_{i=l} = z$  and  $\zeta_{i=h} = 1 - z$ . Set  $\gamma = c/R$  to be the cost ratio of attack, where  $c$  is the unit cost of a single attack. When the attacker’s time preference is described by  $\delta$ , the profit function for a single attacker is

$$\tilde{\Pi}_{A, i} = \int_{t_0}^T e^{-\delta t} \zeta_i \eta_i^{-1} \tilde{\sigma}_i(x_i, \eta_i) dt - \gamma, \quad i \in \{l, h\}. \quad (2)$$

We assume that attackers do not coordinate attacks (or are commissioned by a single attacker) and rewards are claimed on a first-winner-takes-all basis. Attackers are assumed to be drawn from a pool and make one-off entry decisions until marginal cost and marginal benefit are equal and hence  $\tilde{\Pi}_{A,i} = 0$ .

For the targets of such attacks, let  $L > 0$  be an instantaneous value of assets at risk from attack and  $\beta \in \mathbb{R}$  be a subjective discount rate determining the time preferences of all targets. The risk neutral expected loss over the time horizon  $t_0 < t < T$ , is given by

$$\tilde{V}_L = \int_{t_0}^T e^{-\beta t} (z\tilde{\sigma}_l(x_l, \eta_l)L + (1-z)\tilde{\sigma}_h(x_h, \eta_h)L) dt + x_l + x_h. \quad (3)$$

The optimal allocation bundle  $(z^\diamond, x_l^\diamond, x_h^\diamond)$ , when attacking intensity is exogenous, is the simultaneous solution of  $\{\partial\tilde{V}_L/\partial x_l = 0, \partial\tilde{V}_L/\partial x_h = 0, \partial\tilde{V}_L/\partial z = 0\}$ . By construction, if  $\alpha_{i \in \{l,h\}} > 0, \psi_{i \in \{l,h\}} > 0, L > 0, \beta > 0$  and  $z \in (0, 1)$ , a minimum of this function exists. By assumption we set that the optimal allocation must be either  $(x_{i \in \{k,h\}}) \in \mathbb{R}_+$  when  $(\eta_{i \in \{k,h\}}) \in \mathbb{R}_+$  or, if the minimum lies at  $x_{i \in \{l,h\}} < 0$ , then  $x_{i \in \{l,h\}}^\diamond = 0$ . Similarly, we impose the inequality constraint that  $0 \leq z^\diamond \leq 1$ .

Assuming that targets and attackers have positive discount rates the appropriate time horizon,  $T$ , for empirical analysis, maybe determined endogenously. Let  $\lambda$  be an arbitrarily large, but not infinite, number. For a given discount rate,  $\tilde{\theta} = \min(\delta, \beta)$ , by construction

$$\lim_{T \rightarrow \infty} \int_{t_0}^T \tilde{\theta}^{-1} e^{-\theta t} dt = 1.$$

Therefore, the approximation of the time horizon  $\tilde{T}$  covering the  $1 - 1/\lambda$  proportion of the future losses is derived from  $\tilde{T} = \log(\lambda)/\tilde{\theta}$ . In §(3) of this paper, we follow [10] and assume that  $\beta > \delta$  and  $\tilde{T} = \log(\lambda)/\delta$ , such that the interval  $t_0$  to  $\tilde{T}$  covers 90% of the expected present value; that is,  $\lambda = 10$ .

What is important, to the steward, is the overall mass of attacks against systems containing assets under the types  $l$  and  $h$  'storage/operations' areas and this will be influenced by the aggregate behaviour of targets and attackers, rather than the microstructure of individual attack-defence interactions. The more attractive the ecosystem is to attackers, the greater the mass of attacks against its individual components.

**Proposition 1 (Existence of Nash Equilibria Without the Stewart).**

1. (Equilibrium Target Investment) Under the preceding assumptions, when  $\tilde{\sigma}_i = e^{-\psi_i x_i} \eta_i^{\alpha_i}$ , for  $i \in \{l, h\}$ , the Nash equilibrium allocations of  $x_h, x_l$  and  $z$  denoted  $x_h^*, x_l^*$  and  $z^*$  are

$$x_i^* = \frac{\alpha_i}{\psi_i} \log \left( \frac{L\psi_i\psi_j^2(e^{\delta T} - 1)^2}{\gamma\delta\beta(\psi_j + \psi_i)^2} \right) - \frac{\alpha_i\delta T}{\psi_i}, \quad i \in \{l, h\}, j \in \{l, h\}, j \neq i$$

$$z^* = \frac{\psi_l}{\psi_h + \psi_l}. \quad (4)$$

2. (Equilibrium Attacker Intensity) Following from Part 1, above, the Nash equilibrium attacker intensities, denoted  $\eta_l^*$  and  $\eta_h^*$  are

$$\eta_i^* = \left( \frac{\psi_j (e^{\delta T} - 1) e^{-x_i^* \psi_i - \delta T}}{\gamma \delta (\psi_i + \psi_j)} \right)^{\frac{1}{1-\alpha_i}}, \quad i \in \{l, h\}, j \in \{l, h\}, j \neq i, \quad (5)$$

where  $x_{i \in \{l, h\}}^*$ , is the functional forms of the Nash equilibrium given in Part 1 (above).

*Proof.* The proofs of Parts 1 and 2 are given in Appendix A.1. Note that in the multiplicative separably additive form of  $\tilde{\sigma}_{i \in \{l, h\}}$  the Nash equilibrium allocation  $z^*$  is a simple function of  $\psi_{i \in \{l, h\}}$  and when  $\psi_l = \psi_h$  the allocation is equal. If we add the constraint  $x_l + x_h = \tilde{x}$ , where  $\tilde{x}$  is a binding budget constraint, then the attacking effort in each asset area enters the solution for  $z$ .  $\square$

We demonstrate that, in this modelling approach, we do not have to impose an arbitrary constraint on  $x_l + x_h$ , to create conditions similar to the standard results obtained when optimizing under such budget restrictions.

## 2.1 Introducing the Steward

The subject of this paper is resilience, and why a system might not be resilient to security shocks through the choices of the individual components. The first stewardship action we evaluate replicates our previous work (on sustainability [10]) by postulating a Stackelberg policy framework in which the policy-maker stewarding the system sets rules relative to a target level of sustainability. When the steward is fully informed, our model reverts to the mechanism design problem discussed in [10], in which the steward is able to set a mandatory investment bundle (denoted by the lower bar) on the individual targets  $(\bar{x}_l, \bar{x}_h)$ , as well as imposing a specific asset allocation  $\bar{z}$ .

The Nash equilibrium allocations for the  $N_T$  targets assumes no social coordination. Therefore, the Nash equilibrium allocation  $(x_l^*, x_h^*, z^*)$  of defensive effort and corresponding attacking intensities  $(\eta_l^*, \eta_h^*)$  will not necessarily be the first best solution for Pareto efficiency. Let  $(x_l^\dagger, x_h^\dagger, z^\dagger)$  be the Pareto efficient allocations for a given set of model parameters  $(\alpha_{i \in \{l, h\}}, \beta, \gamma, \delta, \lambda, \psi_{i \in \{l, h\}}, L)$ .

A classical efficiency model, with the steward acting as a public policy-maker and imposing  $(\bar{x}_l, \bar{x}_h, \bar{z})$  [10], demonstrates that Pareto efficiency guaranteed only when the subjective discount rate of the steward is equal to  $\beta$ , the common discount rate. Indeed, the analysis in [10] illustrates that, from the subjective viewpoint described by targets' heterogeneous discount rates, the chosen values of  $(\bar{x}_l, \bar{x}_h, \bar{z})$  cannot always be a Pareto efficient allocation,  $(x_l^\dagger, x_h^\dagger, z^\dagger)$ , when  $\beta \neq \bar{\beta}$ . However, there may exist constellations of parameters such the welfare of the individual agents have improved due to the presence of the steward despite their different discount rates. In this study, we do not explore such welfare comparisons.

## 2.2 The Fully Informed $(x_l, x_h, z)$ -setting Steward

Let the steward's discount rate be  $\bar{\beta}$ . A fully informed steward sets a mandatory level of  $(\bar{x}_l, \bar{x}_h, \bar{z})$  by minimizing the following loss function

$$\tilde{V}_P = \int_{t_0}^T e^{-\bar{\beta}t} \left( z \tilde{\sigma}_l(x_l, \eta_l^\diamond) L + (1-z) \tilde{\sigma}_h(x_h, \eta_h^\diamond) L \right) dt + x_l + x_h, \quad (6)$$

where  $\eta_i^\diamond(x_i, z)$  for  $i \in \{l, h\}$  is the solution to

$$\int_{t_0}^T e^{-\delta t} \zeta_i \eta_i^{-1} \tilde{\sigma}_i(x_i, \eta_i) dt = \gamma, \quad i \in \{l, h\}, \quad (7)$$

in terms of  $(x_l, x_h, z)$ . We can see that, by internalizing the attacker reaction curve, the fully informed policy-maker with identical time preferences to the homogenous targets  $\bar{\beta} = \beta$  will set an allocation bundle  $(\bar{x}_l, \bar{x}_h, \bar{z})$ . Moreover, in [10], we show that for the one-dimensional investment case the allocation will be the Pareto efficient allocation from the point of view of both the steward and targets.

In the multi-allocation form of the model, where  $\tilde{\sigma}_i = e^{-\psi_i x_i \eta_i^{\alpha_i}}$ , for  $i \in \{l, h\}$ , proof that  $(\bar{x}_l, \bar{x}_h, \bar{z}) = (x_l^\dagger, x_h^\dagger, z^\dagger)$ , when  $\bar{\beta} = \beta$  for all parameter combinations, is not possible because  $\bar{z}$  does not have a closed form solution (other than in certain special cases; e.g.,  $\alpha_l = \alpha_h = \alpha$ ). One such case is to consider a constraint on weighting aspect of the bundle  $z$  across asset areas of the form:  $z = \psi_h / (\psi_h + \psi_l)$ , the Nash equilibrium allocation. Other constraints on  $z$  can be reasonably justified, as we subsequently demonstrate.

### Proposition 2 (The Fully Informed Steward).

1. (Target investment with steward) When  $\tilde{\sigma}_i = e^{-\psi_i x_i \eta_i^{\alpha_i}}$  and  $\bar{z} = \psi_h / (\psi_h + \psi_l)$ , the stewards optimal investment allocation  $(\bar{x}_l, \bar{x}_h)$  is

$$\begin{aligned} \bar{x}_i &= \frac{1}{\psi_i} \log \left( \psi_j (\psi_i + \psi_j)^{\frac{1}{1-\alpha_j}} \right) + \frac{\alpha_i}{\psi_i} \log \left( \frac{1}{\gamma} \delta (e^{\delta T} - 1) \right) + \\ &\quad \left( \frac{\bar{\beta} T (\alpha_i - 1)}{\psi_i} - \frac{\delta T \alpha_i}{\psi_i} \right) + \frac{(\alpha_i - 1)}{\psi_i} \log \left( \frac{-\bar{\beta} (\alpha_j - 1)}{L \psi_i (e^{\bar{\beta} T} - 1)} \right), \\ &\quad i \in \{l, h\}, j \in \{l, h\}, j \neq i \end{aligned} \quad (8)$$

2. (Attacking intensity) Following from Part 1, above, the attacker intensity  $\eta_{i \in \{l, h\}}$  is

$$\bar{\eta}_i = \left( \frac{\psi_i (e^{\delta T} - 1) e^{-\bar{x}_i \psi_i - \delta T}}{\gamma \delta (\psi_j + \psi_i)} \right)^{\frac{1}{1-\alpha_i}}, \quad i \in \{l, h\}, j \in \{l, h\}, j \neq i \quad (9)$$

where  $\bar{x}_i$  is given in Equation 8.



*Proof.* The proofs of Parts 1 and 2 are given in Appendix A.2. The solution is again subject to an upper bound of  $\eta_i^* < e^{\alpha_i^{-1}x_i\psi_i}$ , for  $i \in \{l, h\}$ . We can compare the solutions in Equations 8 and 9 for the fully informed steward versus those in Equations 4 and 5.  $\square$

**Proposition 3 (The Steward's Improvement).** *If  $\tilde{\sigma}_i = e^{-\psi_i x_i \eta_i^{\alpha_i}}$ , with  $\beta \geq \bar{\beta}$ , and  $\alpha_{i \in \{l, h\}} > 0, \psi_{i \in \{l, h\}} > 0, \gamma > 0, \delta > 0, L > 0$  and the asset allocation is constrained to  $\bar{z} = \psi_h / (\psi_h + \psi_l)$ , then the steward's mandated investment  $\bar{x}_{i \in \{l, h\}}$  is always greater than or equal to the Nash equilibrium investment bundle  $x^*_{i \in \{l, h\}}$ .*

*Proof.* The proof is obtained by substituting the expressions  $\bar{x}_{i \in \{l, h\}}$  and  $x^*_{i \in \{l, h\}}$  in Equations 4 and 8 into the functional form  $\bar{x}_{i \in \{l, h\}} \geq x^*_{i \in \{l, h\}}$  and subject to the constraint  $\beta \geq \bar{\beta}$ . By solving the two inequalities simultaneously for each parameter relative to its own constraint, that is  $\alpha_{i \in \{l, h\}} > 0, \psi_{i \in \{l, h\}} > 0, \gamma > 0, \delta > 0, L > 0$ , by inspection the constraint  $\beta > \bar{\beta}$  is never violated. The complete set of steps of the proof is relatively simple, albeit long algebraic manipulation.  $\square$

A useful by-product of the comparison between Propositions 1 and 2 is that we can define an upper bound on  $\beta \geq \bar{\beta}$  such that the steward does at least as well as the Nash equilibrium even when the steward weights potential near-term losses more than the targets do). Again, this is covered in more detail for the one-dimensional case in [10].

The attacker intensities follow from the functional form of the Nash equilibrium, except with  $\bar{x}_{i \in \{l, h\}}$  replacing  $x^*_{i \in \{l, h\}}$ , as in Equation 9. From the chosen functional form of  $\tilde{\sigma}_{i \in \{l, h\}}, \bar{\eta}_{i \in \{l, h\}}$ , we know that overall loss decreases with increasing  $x_{i \in \{l, h\}}$ , ceteris paribus, and we know, by construction, that  $\bar{x} > x^*$  when we constrain using  $\bar{z} = \psi_h / (\psi_h + \psi_l)$  and  $\beta \geq \bar{\beta}$ .

Following [10], we also consider an non-discounted metric  $\tilde{V}_A$  that measures total cost from attacks and investment. We consider a detailed functional form in § 2.6. If  $x_{i \in \{l, h\}}$  is set by the fully informed steward minimizing the objective function set out in Equation 6 and if  $\tilde{\sigma}_{i \in \{l, h\}}, \bar{\eta}_{i \in \{l, h\}}$ , with  $\bar{z} = \psi_h / (\psi_h + \psi_l)$  and  $\beta \geq \bar{\beta}$ , then  $\tilde{V}_A(\bar{x}_{i \in \{l, h\}})$  will be lower than  $\tilde{V}_A(x^*_{i \in \{l, h\}})$  for all combinations of  $\alpha_{i \in \{l, h\}} > 0, \psi_{i \in \{l, h\}} > 0, \gamma > 0, \delta > 0$  and  $L > 0$ . Although, by construction,  $\tilde{V}_A$  is not an objective function (its minima is unbounded in  $x_{i \in \{l, h\}}$ ), the functional form of  $\tilde{V}_A$  is useful in measuring the effect of the transition from  $x^*_{i \in \{l, h\}}$  to  $\bar{x}_{i \in \{l, h\}}$  free from the subjective discount rates  $\beta$  and  $\bar{\beta}$ .

### 2.3 Reducing the Steward's Abilities

The preceding notion of the steward assumed that it has the ability to impose  $(\bar{x}_l, \bar{x}_h, \bar{z})$  on the targets and thus achieve a lower loss in  $\tilde{V}_P$  than the Nash equilibrium allocation of  $(x_l^*, x_h^*, z^*)$ . This result is useful, if unsurprising. The steward acts as a classic public policy-maker and sets the mechanism so that any attacking externalities are internalised by the targets. A less intuitive fact is

that it is the steward's discount rate  $\bar{\beta}$  that determines if, from the viewpoint of the targets with discount rate  $\beta$ , a Pareto efficient solution has been achieved.

For some parameter combinations of  $\alpha$  and  $\psi$ , with  $\beta \neq \bar{\beta}$ , a natural tension will exist between the targets and the steward. If the steward requires, periodically, say, to have its power to set  $(\bar{x}_l, \bar{x}_h, \bar{z})$  ratified by the targets, then it is likely that  $\bar{\beta} \rightarrow \beta$ . However, if the individual targets have heterogeneous discount rates, then the steward will never be able to attain the Pareto dominant solution unless each target is allowed to state its own discount rate. When this issue occurs, targets may overstate their discount rates — we can consider the security resource allocation to be part of a wider investment bundle for the targets — and their allocation bundles will simply tend back to the Nash equilibrium. We leave extended discussion of this effect to future work.

Moving back to the simplified ex-ante identical targets example, further interesting cases can be analysed by restricting either the action set and/or the information set of the steward. Indeed, these cases present the type of situations where the steward is unable to maintain the resilience of the ecosystem of targets in the presence of shocks to specific parameters (we focus on shocks to the technology parameters  $\alpha_{i \in \{l, h\}}$  and  $\psi_{i \in \{l, h\}}$ ). In the next section, we analyse the cases of the fully informed steward with limited action and, finally, the partially informed steward with limited action.

## 2.4 Full Information with Limited Action: Majority and Minority Cases

First, consider the case in which the steward can observe  $x_{i \in \{l, h\}}$  and  $z$ , but can only impose constraints on  $x_h$  and  $z$ . We designate this the *majority-action-case*; that is, the steward controls the majority of variables affecting the allocation bundle (two variables) and the individual agents control a minority of it (one variable).

A similar case occurs when the the steward can only impose constraints on  $x_h$  and  $x_l$ , but observes  $z$ , the results are intuitively identical. In this case for the targets of attacks,  $x_h$  and  $z$  are now exogenous and their problem reduces to a one-dimensional optimization problem seeking to minimize

$$\tilde{x}_l(z, x_h, \eta_l, \eta_h) = \arg \min_{x_l} \int_{t_0}^T e^{-\beta t} (z \tilde{\sigma}_l(x_l, \eta_l) L + (1 - z) \tilde{\sigma}_h(\bar{x}_h, \eta_h) L) dt + x_l + x_h, \quad (10)$$

where  $\tilde{x}_l(z, x_h, \eta_l, \eta_h)$  is the target's optimal solution for  $x_l$  as a function of the now imposed values of  $x_h$ ,  $z$ , and the attacker intensity choices  $\eta_l$  and  $\eta_h$ . The intuition behind this approach is that the steward sets some collection of rules that identify the allocation  $z$  and then imposes some investment on that allocation  $x_h$ . The optimal bundle of  $(x_h, z)$  from the viewpoint of the steward is denoted  $(\bar{x}_h, \bar{z})$ . The steward therefore solves the other two thirds of the allocation using

the following objective function:

$$\begin{aligned}
(\bar{x}_h, \bar{z}) = & \\
& \arg \min_{x_h, z} \int_{t_0}^T e^{-\bar{\beta}t} \left( (1-z) \tilde{\sigma}_l \left( \tilde{x}_l \left( z, x_h, \eta_l^\diamond, \eta_h^\diamond \right), \eta_l^\diamond \right) L + z \tilde{\sigma}_h \left( x_h, \eta_h^\diamond \right) L \right) dt \\
& + \tilde{x}_l \left( z, x_h, \eta_l^\diamond, \eta_h^\diamond \right) + x_h, \tag{11}
\end{aligned}$$

where  $\eta_{i \in \{l, h\}}^\diamond$  is the solution to the attacker intensities given in Equation 7. As the steward anticipates the reaction of the target into the objective function for  $x_l$ , in this instance, almost all of the steward's objectives in  $(x_l, x_h, z)$  can be achieved. The the steward can impose itself on two out of the three degrees of freedom in the model. We can also see that if  $\bar{\beta} = \beta$  (i.e., when the steward and targets have aligned time preferences), then the steward will achieve a risk profile broadly similar to the case when the steward controls all of the degrees of freedom  $(x_l, x_h, z)$ . Whilst the steward can attain its desired risk expenditure trade-off, it can do so only at a lower level of efficiency (in terms of total initial cost  $x_l + x_h$ ) than if the steward controls  $(x_l, x_h, z)$ . Unless an arbitrary upper bound is placed on  $x_h + x_l$ , the steward can achieve a global minimum, for any given combination of  $\alpha_{i \in \{l, h\}}$  and  $\psi_{i \in \{l, h\}}$ , by imposing a shift of assets (if necessary) into the high security domain. In the extreme case, in which  $\bar{z} \rightarrow 1$ , the steward has control over all assets and sets an unbounded investment in protection of  $\bar{x}_h$  as an essentially one-dimensional optimization problem.

Reducing the steward's action space to only one of the three allocation variables (which we call the minority action case) provides a far greater limitation to its action space and substantially impairs the steward's ability to internalize the attacker externalities and adjust the total level of risk in response to a change in  $\alpha_{i \in \{l, h\}}$  or  $\psi_{i \in \{l, h\}}$ . However, the circumstances in which a steward would be able to observe behaviour, but have no direct influence on it, violate one of the presumed key roles of the steward in the ecosystem leave the motivation and analysis of this fully informed, but substantively limited, steward to future work.

## 2.5 The Partially Informed Steward with Limited Action: Minority Case

We skip case of the fully informed steward with limited action and move directly to a partially informed steward with minority action. This, in theory at least, is the most interesting case as it illustrates both the limitations of the steward's actions in response to changes in  $\alpha_{i \in \{l, h\}}$  or  $\psi_{i \in \{l, h\}}$  and also that, with limited information, the presence of the steward can in fact lead to a worse global outcome than the Nash equilibrium .

Let the steward observe and enforce only  $x_h$ . The steward can observe and internalize the externality in  $\eta_h$ , but cannot observe or enforce  $z$  or  $x_l$ . The

targets then choose the investment and allocation bundle  $(x_l, z)$  following

$$(\tilde{x}_l, \tilde{z}; x_h, \eta_l, \eta_h) = \arg \min_{x_l, z} \int_{t_0}^T e^{-\beta t} (\tilde{z} \tilde{\sigma}_l(x_l, \eta_l) L + (1 - \tilde{z}) \tilde{\sigma}_h(\tilde{x}_h, \eta_h) L) dt + x_l + x_h. \quad (12)$$

The steward now solves the following minority optimization, with the steward's given information set:

$$\bar{x}_h(\eta_h) = \arg \min_{x_h} \int_{t_0}^T e^{-\beta t} \left( \hat{L} \tilde{\sigma}_h(\bar{x}_h, \eta_h^\diamond) \right) dt + x_h, \quad (13)$$

where  $\eta_h^\diamond$  is the solution to the attacker entry problem from Equation 2, but only for the  $h$  asset class.

From the steward's point of view this is now

$$\int_{t_0}^T e^{-\delta t} \tilde{\zeta}_h \eta_h^{-1} \tilde{\sigma}_h(x_h, \eta_h) dt = \gamma. \quad (14)$$

Note that the steward now takes for given  $\hat{L}$  as the value of losses; this is because the steward can no longer identify  $zL$  and  $(1 - z)L$ , the steward is simply given  $\hat{L}$  by the targets at an a-priori stage and is assumed to be exogenous. Similarly, whilst  $\tilde{\zeta}_h$  is equal to  $z$  from the viewpoint of attackers and targets, it is simply a parameter unrelated to the overall asset allocation of the targets from the point of view of the steward. The steward is now unwittingly, not a Stackelberg policy maker, but in a Nash equilibrium with the targets and attackers.

The attackers are also solving their entry and exit decision for assets in allocation  $l$ , following  $\int_{t_0}^T e^{-\delta t} \tilde{\zeta}_l \eta_l^{-1} \tilde{\sigma}_l(x_l, \eta_l) dt = \gamma$ . This is unobserved by the steward, but is accounted for as part of a Nash equilibrium by the targets. For tractability, we assume that, from the viewpoint of the attackers,  $\tilde{\zeta}_h$  is set exogenously and at a fixed ratio to  $\hat{L}$ . We are interested in the reaction of targets setting  $x_l$  and attackers choosing  $\eta_l$ , in order to demonstrate the natural limits that appear in the game and to analyse this case, we assume without loss of generality that  $\hat{L}$  is exogenous by construction and  $\tilde{\zeta}_h$  is already set in a pre-optimization between the attackers and the steward.

**Proposition 4 (Attackers and Steward).**

1. (Asset Class  $h$ ) If  $\tilde{\sigma}_i = e^{-\psi_i x_i} \eta_i^{\alpha_i}$ , for  $i \in \{l, h\}$ , the steward's objective function is as stated in Equation 13, and the attacker dynamics are as given in Equation 14, then the steward's optimal mandated investment allocation is

$$\bar{x}_h = \frac{1 - \alpha_h}{\psi_h} \log(\hat{L}(e^{\beta T} - 1)\psi_h) - \frac{\alpha_h}{\psi_h} \log(\gamma \delta (\tilde{\zeta}_h e^{\delta T} - \tilde{\zeta}_h)) + \frac{1}{\psi_h} (\log(\bar{\beta} \alpha_h - \bar{\beta})(1 - \alpha_h) + \alpha_h T(\bar{\beta} - \delta) - \bar{\beta} T). \quad (15)$$

Following from the steward's choice, the attacker intensity, given the steward's actions  $\bar{\eta}_h$ , is given by

$$\bar{\eta}_h = \left( \frac{\tilde{\zeta} (e^{\delta T} - 1) e^{-\bar{x}_h \psi_h - \delta T}}{\gamma \delta} \right)^{\frac{1}{1-\alpha_h}}, \quad (16)$$

where  $\bar{x}_h$  is as defined in Equation 15.

2. (Asset Class  $l$ ) We now consider the targets' and attackers' new equilibrium: if  $\tilde{\sigma}_i = e^{-\psi_i x_i} \eta_i^{\alpha_i}$ , for  $i \in \{l, h\}$ , and the targets' objective is as specified in Equation 12, then the equilibrium allocation bundle  $x_l, z$  will be

$$x_l^\ddagger = -\frac{1}{\psi_l} \log(\bar{\eta}_h^{\alpha_h}) + \frac{\alpha_l}{\psi_l} (\log(\bar{\eta}_h^{\alpha_h}) + \log(\beta(e^{\delta T} - 1)\bar{\eta}_h^{-\alpha_h}) - \quad (17)$$

$$\log(\gamma \delta \psi_l (e^{\beta T} - 1)L) + \beta T - \delta T) + \bar{x}_h \psi_h$$

$$z^\ddagger = \frac{\beta \bar{\eta}_h^{-\alpha_h} e^{\bar{x}_h \psi_h + \beta T}}{L \psi_l (e^{\beta T} - 1)}, \quad (18)$$

and the attacker intensity  $\eta_l$  is given by

$$\eta_l^\ddagger = \left( \frac{z (e^{\delta T} - 1) e^{-x_l^\ddagger \psi_l - \delta T}}{\gamma \delta} \right)^{\frac{1}{1-\alpha_l}}. \quad (19)$$

*Proof.* The proof of Parts 1 and 2 are given in Appendix A.3. Note that the steward's choice is effectively determined by three variables  $\hat{L}$ ,  $\tilde{\zeta}_h$ , and  $\bar{\beta}$ . We assume that these are, a priori, in the steward's information set. It is worth reiterating that, at this stage, decisions regarding  $z$ ,  $x_l$ , and  $\eta_l$  are, by construction, not included in this optimization. However, we do not have to impose these restrictions, as the derivative with respect to  $x_h$  of the steward's objective function, given the multiplicative form of  $\tilde{\sigma}_i = e^{-\psi_i x_i} \eta_i^{\alpha_i}$ , for  $i \in \{l, h\}$ , does not include  $x_l$  and  $\eta_l$ . So, the only implicitly restricted information is replaced by  $\hat{L}$ ,  $\tilde{\zeta}_h$ .  $\square$

Note that we use the  $\ddagger$  to denote this new equilibrium for the targets as it is not strictly a Nash equilibrium solution, but rather is Bayes-Nash equilibrium, in which the steward has prior values for  $\hat{L}$  and  $\tilde{\zeta}$ . See [7] for an explanation of the difference between Nash and Bayes-Nash equilibria.

## 2.6 Measuring Resilience

Measuring the impact of technological shocks to  $\alpha_{i, i \in \{l, h\}}$  and  $\psi_{i, i \in \{l, h\}}$  and economic shocks to  $\bar{\beta}$ ,  $\beta$ ,  $\delta$ ,  $L$ , and  $\gamma$  is a challenging task and requires the creation of an arbitrary metric. In this paper, we combine the equilibrium values of  $x_{i \in \{l, h\}}$ ,  $z$ , and  $\eta_{i \in \{l, h\}}$  using a total non-discounted loss function for the risk

component only. This is given as follows:

$$\tilde{V}_A(\tilde{v}, \tilde{u}) = \int_{t_0}^{\tilde{T}} \tilde{z} \tilde{\sigma}_l(\tilde{x}_l, \tilde{\eta}_l) L + (1 - \tilde{z}) \tilde{\sigma}_h(\tilde{x}_h, \eta_h) L dt \quad (20)$$

$$\tilde{v} = (\tilde{z}, \tilde{x}_{i \in \{l, h\}}, \tilde{\eta}_{i \in \{l, h\}}) \quad (21)$$

$$\tilde{u} = (\alpha_{i, i \in \{l, h\}}, \psi_{i, i \in \{l, h\}}), \quad (22)$$

where  $\tilde{T} = -\log(\lambda)/\theta$  and  $\theta = \min(\bar{\beta}, \beta, \delta)$ , for an arbitrary  $\lambda$  tending to zero. By construction, Equation 20 gives an undiscounted loss function, so that the value of the critical parameter  $\tilde{T}$ , which represents the step-size of the periods considered in the model (cf. Figure 2, for a multi-period sustainability model), is finite.  $\tilde{v}$  is the collection of choice variables under the various stewardship options.  $\tilde{u}$  is the collection of parameters that are subject to the technology shocks under consideration.

For a single period, resilience will be measured by a response function to shocks to the parameters  $\tilde{u}$ . Our choice of response function to technology shocks allows for shocks across the set of parameters  $\tilde{u}$  either simultaneously or individually. It is given by the numerical evaluation of the following ordinary differential equation:

$$\tilde{I}(\tilde{u}) = \int_{t_0}^{\tilde{T}} \frac{\partial \tilde{z}}{\partial \tilde{u}} \tilde{\sigma}_l \left( \frac{\partial \tilde{x}_l}{\partial \tilde{u}}, \frac{\partial \tilde{\eta}_l}{\partial \tilde{u}} \right) L + \frac{\partial (1 - \tilde{z})}{\partial \tilde{u}} \tilde{\sigma}_h \left( \frac{\partial \tilde{x}}{\partial \tilde{u}}, \frac{\partial \tilde{\eta}}{\partial \tilde{u}} \right) L dt, \quad (23)$$

$$\tilde{u} = \{\alpha_{i \in \{l, h\}}, \psi_{i \in \{l, h\}}\},$$

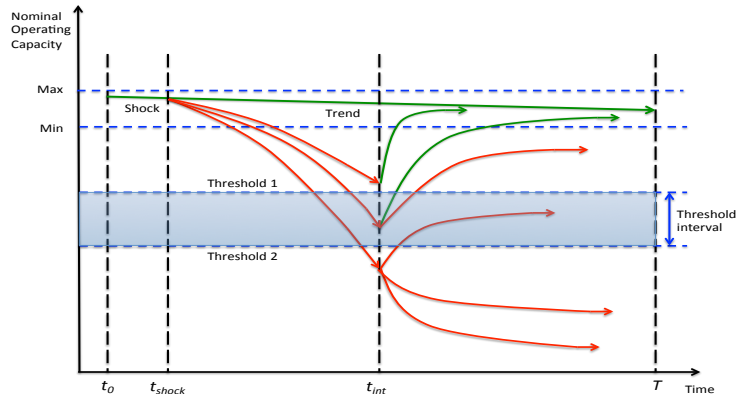
where each case has a set of functional forms for  $\tilde{z}$ ,  $\tilde{x}_{i \in \{l, h\}}$  and  $\eta_{i \in \{l, h\}}$ . We have denoted the three cases as follows:  $\tilde{v} = v^*$  and  $\tilde{u} = u^*$  for the Nash equilibrium,  $\tilde{v} = \bar{v}$  and  $\tilde{u} = \bar{u}$  for the fully informed steward and  $\tilde{v} = \bar{v}^\ddagger$  and  $\tilde{u} = \bar{u}^\ddagger$  for the partially informed steward with minority action case.

We are interested in establishing the thresholds, illustrated in Figure 4, which describe levels of system operating capacity, as measured by loss, for differing degrees of steward's effectiveness. We attempt to establish whether the system restores, through co-ordinated investment, to the target zone or not.

In our model, these thresholds reveal themselves as discontinuities, relative to shock size, in the solutions to Equation 23, below. Such discontinuities can be seen in our simulations as the asymptotes in Figures 6 and 8, for fully and partially informed stewards.

Let  $\tilde{V}_A(v^*, u^*)$  and  $\tilde{I}(u^*)$  be, respectively, the total non-discounted loss for the risk component under the Nash equilibrium and the corresponding collection of response functions. Similarly let  $\tilde{V}_A(\bar{v}, \bar{u})$ ,  $\tilde{I}(\bar{u})$  and  $\tilde{V}_A(\bar{v}^\ddagger, \bar{u}^\ddagger)$ ,  $\tilde{I}(\bar{u}^\ddagger)$  be, respectively, the same pair of functions and collection of functions for the fully informed steward and the partially informed steward with minority action cases.

We can measure the effectiveness of the steward by comparing  $\tilde{V}_A(v^*, u^*)$  to  $\tilde{V}_A(\bar{v}, \bar{u})$ . We can also evaluate the erosion in risk reduction caused by re-



**Fig. 4.** Resilience with an incompletely informed steward.

stricting the stewards information set and action space by pairwise evaluation of  $\tilde{V}_A(v^*, u^*)$  and  $\tilde{V}_A(\bar{v}, \bar{u})$  with  $\tilde{V}_A(\bar{v}^\ddagger, \bar{u}^\ddagger)$ .

To examine the impact of shocks and measure resilience we compare the response functions  $\tilde{I}(u^*)$  and  $I(\bar{u})$  to evaluate the impact of the fully informed steward. Finally, we can compare the resilience of the system when the stewards information set is restricted by comparing  $\tilde{I}(u^*)$  and  $I(\bar{u})$  to  $I(\bar{u}^\ddagger)$ , for varying sizes of shocks in  $\bar{u}$ . In particular, we focus on  $\alpha_{i \in \{l, h\}}$ .

### 3 Application to ICS, SCADA, and Corporate Networks

Industrial control systems (ICS) are ubiquitous in most large industrial firms and related organisations. A further common type of ICS are Supervisory Control and Data Acquisition (SCADA) systems. These systems are designed to automatically or semi-automatically control industrial processes. Examples of such systems can be found in petroleum exploration and processing, gas distribution, bulk electricity transmission, various parts of the nuclear industry and most manufacturing processes. Similar, or identical types of systems may be found in defensive applications, such as automatic air defence systems. ICS/SCADA systems are often very complex and deal with a large number of different types of sensors and actuators affecting the various aspects of the system in question. ICS/SCADA systems and the security of ICS/SCADA systems is not a new topic however, when many of the ICS/SCADA systems were first installed they were viewed as standalone assets and as such the major security concern was physical access to the control system or by physically tapping directly into the data acquisition sensors and/or the control communications to actuators. For our purposes, this distinction between ICS and SCADA is not critical and

we refer generically to ICS/SCADA as a single type of assets within a target organization.

Our main question centres on whether a firm would seek to adjust its declared mix of ICS/SCADA and corporate information assets (we explicitly do not include physical assets in this example) to avoid costly regulation. We will assume that there exists some legacy regulation of certain types of ICS/SCADA systems and that firms can choose to replace some or all of the information architecture of these systems with analogous technologies run on an unregulated corporate network. In terms of the model presented here, we have following set-up:

	investments	allocation	risk-reduction rate	attacker elasticity
ICS/SCADA	$x_h$	$1 - z$	$\psi_h$	$\alpha_h$
Corporate	$x_l$	$z$	$\psi_l$	$\alpha_l$

In this paper, we run simulations for this model in which we shock the  $\alpha_l$ ; that is, the elasticity of attacker intensity against assets in the corporate network. Here we are modelling the situation in which we assume that the primary source of vulnerability is associated with the corporate network because of its more direct exposure to the Internet, with all of the associated vulnerabilities. As ICS/SCADA systems are increasingly interconnected with corporate systems, these vulnerabilities potentially affect core CNI systems; that is, the high-value ( $h$ ) assets. Clearly, Equation 23 allows for a wider variety of experiments, which would allow us to explore different assumptions about the sources of vulnerabilities.

In the US, 1,900 bulk power system operators are regulated by The North American Electric Reliability Corporation (NERC), a not-for-profit organization with the role of coordinating the individual operators. This regulated ecosystem — of interconnected organizations — provides us with some indicative parameters for our experiments. Each operator will have a ICS/SCADA system that manages the bulk electricity transmission in their area. This will be a network of communications that monitor the physical network of power cables, transformers, and substations.

In addition to the ICS/SCADA assets, the various operators have corporate networks that provide on-going information services for the normal business activities for each operator. The corporate network has many of the same features as the ICS/SCADA system and there are elements of substitutability between the two. For instance, an operator could phase out using expensive fibre optic cables to communicate between ICS/SCADA systems and substations and replace this with a IP or 3G type communications.

A successful penetration of a corporate network that is integrated with an ICS/SCADA now provides attackers with a potentially more effective means of attacking the ICS/ SCADA system. The attacker can sit and learn the systems properties via sampling and observation of the ICS/SCADA systems normal operation and then use this information to either provide a priori information to improve the chance of success of a physical attack or actually attack the ICS/SCADA system directly through the corporate network.



As a community of targets, systematic underinvestment across all targets leads to increased attacking intensity and this provides a negative externality that requires coordination across targets in order to internalize this cost. We will illustrate three cases for this example, first where targets are unregulated and choose investment using the Nash equilibrium approach in § 2. We will then demonstrate the improvement that can be achieved by the fully informed steward. Finally, we will illustrate the deterioration in security when targets can shift assets from the oversight of the steward and the steward can no longer mandate investment. In each case, we illustrate the change in total risk with shocks to attacker elasticities and why targets may find it attractive to move assets from a regulated to an unregulated environment.

### 3.1 An Example Simulation

This simulation is designed to provide an overview of the intuition of our model and is not supposed to provide specific quantification for our proposed application. However, we have tried to stay close to real data when possible.

Let us assume that targets have a discount rate of 20% per annum ( $\beta = \log(6/5)$  continuous growth rate), in this case when  $\lambda = 10$ , the target time overall horizon is  $T = 12.3$  years. This appears to be a reasonable assumption for the amortization of information assets within a firm see, for example, the survey in [4]. For electricity transmission in the United States, the difference between physical and information assets can be found in [12, 14].

We assume that the societal discount rate used by the steward is much lower and ranges from  $\bar{\beta} \rightarrow 0$  to  $\bar{\beta} \rightarrow 1/10$ . In [10], we outline the various debates on the appropriate social discount factor to be applied in public policy scenarios. For certain areas of public policy debate such as climate change discount rates approaching zero a common for certain economic arguments relating to low carbon policies. For information stewardship the requirement is not so acute but significant differences between firm discount rates and societal discount rates remain.

For our starting numerical example, we assume that  $\psi_l = \psi_h$ ; that is, the relative marginal risk reduction from investment in both asset classes is identical and fixed we assume that it is  $1/100$ ,  $1/10$  and  $1/2$ , to represent low, medium and high effectiveness bands. This is a more difficult assumption to justify as there is very little literature on the efficacy of investment in security in this area, therefore our simulation covers a wide range of reasonable bands.

We arbitrarily fix  $L = \$1\text{M}$ , as an example, and divide all losses by  $L$  to give a per-dollar-at-risk measure.  $\hat{L}$  is assumed to be half  $L$ . Starting from the Nash equilibrium assumption, if  $\psi_l = \psi_h$ , it follows that  $\tilde{\zeta} = 1/2$ .

We set the attackers' discount rate to be  $\delta = \log(11/10)$ , or a 10% discrete rate of return. From the viewpoint of attackers, the discount rate is analogous to an investment, as opposed to depreciation and amortization from the viewpoint of the targets. The most difficult parameter to set in the simulation is  $\gamma$ , as almost no data exists on the cost per attack to reward ratio. When  $\gamma \rightarrow 0$ , the cost per attack divided by reward indicates that either the rewards are very high or that

the cost per attack is very low. When  $\gamma = 0$ , attacking intensity is infinite. This has not been observed, therefore we stick to finite values of  $\gamma = 1/10$  or a 10% cost-reward ratio. The shock of interest is that to the elasticity of attack  $\alpha_{i \in \{l, h\}}$  and, in particular, shocks to  $\alpha_l$ .

Figures 5 and 6 illustrate the differing effects of shocks to the attacker elasticities, in the presence of a fully informed steward. Recall that by increasing the elasticity the attacker chances of success increase substantially, *ceteris paribus*.

However, the ecosystem will react to shocks to this elasticity, for the Nash equilibrium in the absence of a steward this will simply be an adjustment to the ratios illustrated in Proposition 1. However, for the fully informed steward there is a reactive coordinating entity, balancing current period investment with future uncertain losses. As shocks to the attacking elasticity  $\alpha_l$  increase, the steward utilizes this collective component to reduce the attacking intensity (rather than keeping the risk down by defensive effort  $x_l$ ). The derivative of  $\partial \tilde{\eta}_l / \partial \bar{x}_l$  is now more important than  $\partial \tilde{\sigma}_l / \partial \bar{x}_l$ , where  $\eta_l$  is constant. The steward therefore finds an optimum by driving away all the attackers (as even small numbers are now very effective).

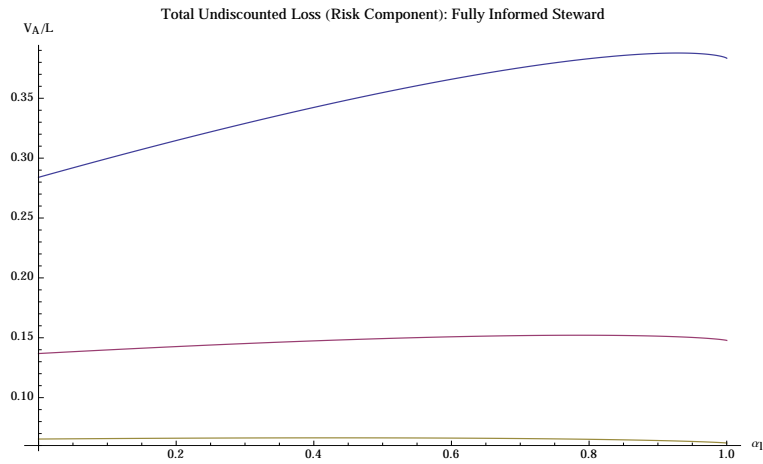
We see that, for all values of  $\psi$ , the fully informed steward provides a lower total non-discounted loss than the Nash equilibrium. This illustrates the beneficial effect of the steward. However, with larger values of  $\psi$ , the absolute effect decreases. The major benefit of the steward is in suppressing and adjusting the ecosystem to shocks and this effect is demonstrable for all three values of  $\psi$ .

Finally, we move to the partially informed steward with minority action, the total non-discounted loss  $\tilde{V}_A$  and response function  $\tilde{I}$  for shocks in  $\alpha_l$  are plotted in Figures 7 and 8. In this case, the pattern is similar to the Nash equilibrium for small shocks. The targets, however have costly regulation in the  $h$  asset class and are under investing in the  $l$  asset class. Unfortunately, in this case there is a discontinuity at  $\alpha_l = 1$ , so the total loss spikes prior to the shift in assets from  $l$  to  $h$ . This is a *de facto* boundary, as illustrated in Figure 4. We can see that before the steward can regulate the assets, the total risk will traverse the discontinuity, before the steward can actually manage the majority of assets that the targets have not declared. Here, we can see a case of an ecosystem that is not resilient and lies within the feasible boundaries of our example parameter sets.

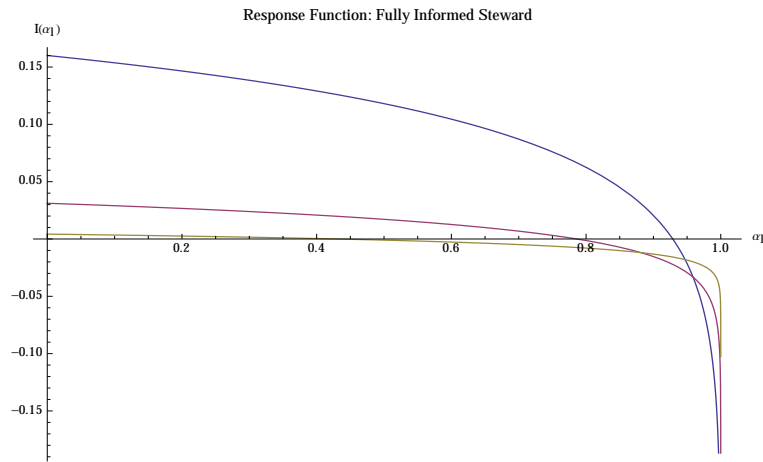
### 3.2 Robustness of the Modelling Assumptions

The various forms of the model that we have proposed assume that targets are *ex-ante* identical. This is, of course, a simplifying assumption to lend tractability to the derivation and illustration of the specific effects that we are attempting to identify. However, this assumption is not as limiting as might be suspected.

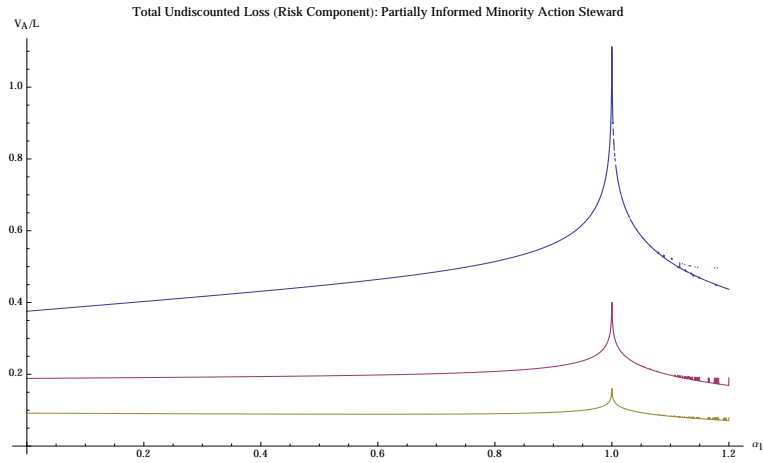
The issue with the heterogeneity of the types of target — in terms of vulnerability or magnitude of loss — is that once we assume a steward in the role of a policy-maker determining mandatory investments, this steward would necessarily have to identify each target's Pareto efficient investment. For a large



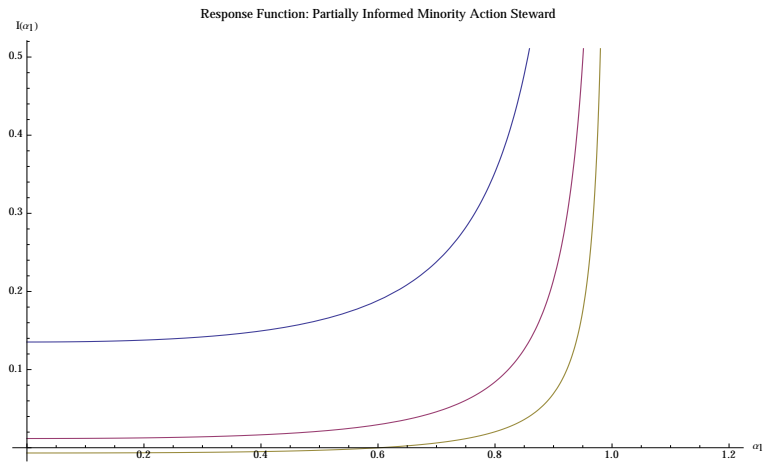
**Fig. 5.** Steward's total non-discounted loss function,  $\tilde{V}_A$ , as a function of  $\alpha_l$ . An important point to note is that this does not include the deterministic up-front investment, so this curve can actually slope downwards, even with increasing  $\alpha_l$ . The upper curve represents  $\psi_l = \psi_h = \psi = 0.01$ , the middle curve  $\psi = 0.1$  and the lower curve is  $\psi = 0.5$ . These values of  $\psi$  represent, respectively, low, medium, and high rates of risk reduction for additional investment.



**Fig. 6.** Steward's response function,  $\tilde{I}(\alpha_l)$ , as a function of an increasing shock in  $\alpha_l$ , the abscissa values. Note that the steward now takes a positive action and seeks to manage the direction of the shock, as  $\alpha_l$  becomes very large, the steward tolerates almost no attacking intensity and this effect is illustrated by the change in sign of the response.



**Fig. 7.** Partially informed steward with minority action total non-discounted loss function  $\tilde{V}_A$  as a function of  $\alpha_l$ . In this case, the targets maintain assets in the increasingly risky  $l$  class to avoid ostly regulation in  $h$ , however a discontinuity exists at  $\alpha_l$  causing the loss function to spike before the assets are shifted back to the regulated domain.



**Fig. 8.** Partially informed steward with minority action response function  $\tilde{I}(\alpha_l)$  as a function of an increasing shock in  $\alpha_l$ , the abscissa values. Note that after a shock of  $\alpha_l \rightarrow 1$ , the function  $\tilde{I}(\alpha_l)$  is not defined.

cross section of targets, this could potentially be a costly information-gathering exercise.

Targets have the incentive to under-disclose their characteristics (e.g., because of budget pressures) and the remediation action of the steward is therefore rendered ineffective. A standard approach to this is contingent audit; see, for example, early research in this area in [11, 2] and later work in [6, 5]. Targets are asked to declare their characteristics — in terms of vulnerability and magnitude of loss — by the steward. In the event of an incident, there is a chance of audit (with known likelihood) and a large penalty (necessarily large enough for incentive compatibility) for incorrect prior identification to the steward. If the target has correctly identified their characteristics then no fine is levied. For the types of model proposed in [13, 10], this approach would allow the steward to coordinate and mandate investment allocations with targets declaring their own vulnerability and loss characteristics. The allocation would therefore be Pareto efficient from the viewpoint of the steward. However, the allocation will not necessarily be Pareto efficient from the viewpoint of the target as the steward and target time preferences may be divergent. This is further exaggerated when the targets have the ability to hide assets from the steward. If the steward's discount rate is very low relative to the targets' rates, then, under certain cases of the model, targets will move their assets to the class labelled  $l$ , by decreasing  $z$  substantially towards zero. This leaves very few assets in class  $h$  regulated by the steward.

When shocks — to the elasticity of the technology of attack in class  $l$  denoted  $\alpha_l$ , say — result in a higher level of viable attacking intensity in equilibrium, targets can either choose to shift their assets to  $h$  by decreasing  $z$  or try to cope with the increasing attacks in  $l$ . Unfortunately, the game between attackers and targets in  $l$  results an equilibrium with externalities. Moreover, for certain versions of the model, the total risk when the steward takes action without observing  $x_l$  and  $z$  may be substantially higher (by orders of magnitude) than if the targets and attackers achieved a Nash equilibrium in the absence of the steward. We have discussed this case in §3.

Several rationales can be put forward to explain why the common knowledge assumption of  $z$  and  $x_l$  might not be shared with the steward by the targets. First, if  $\beta$  is much larger than  $\bar{\beta}$ , then the targets do not share the sustainability objectives of the steward, defined in terms of their time preferences (the targets are far shorter term than the policy-maker), therefore the targets may make a strategic choice, in an initial sub-game, to hide  $z$  and  $x_l$  from the steward. Second, an alternative explanation, that does not require another mechanism to explain it, is that the targets and steward initially entered into a Stackelberg arrangement that is binding to the steward (to accomplish some sustainability target and internalize externalities in  $x_h$ ). The steward sets  $x_h$  within the framework of the original agreement and this optimization rule continues through the life of the ecosystem, even when potentially new assets  $x_l$  exist. Indeed, the steward may simply not have sufficient information processing power to supervise all assets and then to cover them under relevant tort law liability conditions for

the targets self-revelation approach to work. If there are a very large number of targets with highly diverse information assets, then the full audit may not be possible. Clearly, the model assumes the types of organization in  $x_l$  are ex-ante homogeneous.

One can postulate a set of regulations (in the form of fixed rules) designed by the steward and requiring the disclosure of targets' assets such that the investment  $x_h$  internalizes attacker externalities across targets (on the assumption that this is the complete set of assets). However, after a time, new assets not covered by the rules appear, or methods that allow targets to de-recognize these assets from the steward may exist.

## 4 Summary

This paper will make grim reading for any governmental, supra-governmental agency or firm that needs to act in a stewardship capacity over a complex information ecosystem. We illustrate two contrasting issues that complicate the management of this type of ecosystem. First, for almost all conceivable target–attacker interactions the presence of a steward is beneficial to overall risk reduction, by acting as a social coordinator and mandating investment that internalises externalities. Second, it is unlikely, however, that the time preferences of the steward, acting on behalf of society, and the targets will be aligned and as such the targets may not have the correct incentive to reveal their true type to the steward. In our framework this is in the form of hiding assets in an alternative unregulated asset class.

If the steward is able to observe these assets and mandate the majority of the investment bundle then the steward can still perform a beneficial role. However, when the steward acts on minority information and has limited action, the effect can be far worse than the Nash equilibrium when the steward is not present. Targets, maybe incentivized to store assets in increasingly insecure areas and this can substantially degrade the resilience of the ecosystem.

We have also provided a short example of this model using parameters designed to approximate the choice between holding information assets in a regulated ICS/SCADA system versus redeployment to a standard corporate information network. We demonstrate that a catastrophic scenario predicted by the model solutions under certain parameter configurations is possible for the domain of shocks assumed choices in this example.

Our major conclusions are also backed by qualitative analysis of the types of contracts and regulations needed to ensure that the stewards information set is sufficient to maintain the information ecosystem. It should be noted that the types of regulatory structures needed, transparent mandatory information sharing, audit in the event of an incident and penalties that compound when prior information sharing has proven to be false, are not commonly observed in the regulation of standard enterprise networks. However, for firms with assets of significant importance or interaction with to the critical infrastructure of nations and groups of nations this type of regulation may need to be commonplace.

This could include supply chains to infrastructure providers and cloud service providers for very large groups of firms and individuals.

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## A Nash Equilibrium

### A.1 Proof of Proposition 1

**Part 1: Equilibrium Target Investment** Let  $\tilde{\sigma}_{i \in \{l, h\}} : \mathbb{R}_+ \rightarrow [0, 1]$ . Evaluating the non-stochastic integral of loss over  $t_0 = 0$  to  $T$  we find an analytic form the loss function,

$$V_L = \frac{1}{\beta} L \left( e^{\beta T} - 1 \right) e^{-x_h \psi_h - x_l \psi_l - \beta T} \left( z e^{x_h \psi_h} \eta_l^{\alpha_l} - (z - 1) \eta_h^{\alpha_h} e^{x_l \psi_l} \right) + x_h + x_l. \quad (24)$$

Differentiating with respect to  $\tilde{x}_l$ ,  $\tilde{x}_h$ , and  $\tilde{z}$  yields

$$\frac{\delta \tilde{V}_L}{\delta x_l} = 1 - \frac{1}{\beta} L z \psi_l \left( e^{\beta T} - 1 \right) \eta_l^{\alpha_l} e^{-x_l \psi_l - T \beta} \quad (25)$$

$$\frac{\delta \tilde{V}_L}{\delta x_h} = 1 - \frac{1}{\beta} L (1 - z) \psi_h \left( e^{\beta T} - 1 \right) \eta_h^{\alpha_h} e^{-x_h \psi_h - T \beta} \quad (26)$$

$$\frac{\delta \tilde{V}_L}{\delta z} = \frac{1}{\beta} L \left( e^{\beta T} - 1 \right) e^{-x_h \psi_h - x_l \psi_l - T \beta} \left( e^{x_h \psi_h} \eta_l^{\alpha_l} - \eta_h^{\alpha_h} e^{x_l \psi_l} \right). \quad (27)$$

Setting  $\delta \tilde{V}_L / \delta x_l = 0$ ,  $\delta \tilde{V}_L / \delta x_h = 0$  and  $\delta \tilde{V}_L / \delta z = 0$ , and solving simultaneously, we derive the unconstrained optimal allocation  $(x_l^\diamond, x_h^\diamond, z^\diamond)$ . When attacking intensity  $(\eta_l, \eta_h)$  is exogenous, this is analytically derived as

$$x_l^\diamond(\eta_l) = \frac{1}{\psi_l} \log \left( \frac{(\lambda - 1) L \psi_h \psi_l \eta_l^{\alpha_l}}{\beta \lambda (\psi_h + \psi_l)} \right) \quad (28)$$

$$x_h^\diamond(\eta_h) = \frac{1}{\psi_h} \log \left( \frac{(\lambda - 1) L \psi_h \psi_l \eta_h^{\alpha_h}}{\beta \lambda (\psi_h + \psi_l)} \right) \quad (29)$$

$$z^\diamond = \frac{\psi_l}{\psi_h + \psi_l}. \quad (30)$$

Note that  $z^\diamond$  is a simple ratio of  $\psi_h$  and  $\psi_l$ . In this model, we apply no total budget constraint on  $x_h$  and  $x_l$ ; that is,  $x_h + x_l = x$ , so no Lagrange multiplier needs to be added at this stage.

**Part 2: Equilibrium Attacker Intensity** Following from the target decision-making process, we derive the attacker intensity function. Attackers enter the market for attacks in each asset class until they break even. When  $\Pi_l = 0$  and  $\Pi_h = 0$ , we assume that attackers are randomly assigned targets, with identical probability  $1/N_T$  for each attack, and that the first successful attacker wins the reward  $R$ .

Let  $\gamma = c/R$ , the cost of attack to reward. When  $\tilde{\sigma}_{i \in \{l, h\}} : \mathbb{R}_+ \rightarrow [0, 1]$ , the profit functions for the attacker are as follows:

$$\Pi_l = \frac{1}{\delta} z \lambda^{-\frac{\delta}{\beta}} \left( \lambda^{\delta/\beta} - 1 \right) \eta_l^{\alpha_l - 1} e^{-x_l \psi_l} - \gamma \quad (31)$$



$$\Pi_h = \frac{1}{\delta}(1-z) \left( \lambda^{\delta/\beta} - 1 \right) \lambda^{-\frac{\delta}{\beta}} \eta_h^{\alpha_h-1} e^{-x_h \psi_h} - \gamma. \quad (32)$$

Solving each function for the break-even attacking intensities  $\eta_l^\diamond(x_l)$  and  $\eta_h^\diamond(x_h)$ , we compute the aggregate attacker reaction functions:

$$\eta_l^\diamond(x_l) = \left( \frac{z \lambda^{-\frac{\delta}{\beta}} (\lambda^{\delta/\beta} - 1) e^{-x_l \psi_l}}{\gamma \delta} \right)^{\frac{1}{1-\alpha_l}} \quad (33)$$

$$\eta_h^\diamond(x_h) = \left( \frac{(1-z) \lambda^{-\frac{\delta}{\beta}} (\lambda^{\delta/\beta} - 1) e^{-x_h \psi_h}}{\gamma \delta} \right)^{\frac{1}{1-\alpha_h}}. \quad (34)$$

The simultaneous Nash equilibrium is the best reply of the target to the best reply of the attacker (and vice versa), which is the simultaneous solution of  $\{x_l^\diamond, x_h^\diamond, z^\diamond, \eta_l^\diamond, \eta_h^\diamond\}$ .

Setting the Nash equilibrium defensive allocation (targets) and attacking intensity (attacker) as  $\{x_l^*, x_h^*, z^*, \eta_l^*, \eta_h^*\}$ , we obtain

$$\begin{aligned} x_l^* &= \frac{\alpha_l}{\psi_l} \left( -\log(\gamma \delta L \psi_h \psi_l (e^{\beta T} - 1)) + \log(\beta \psi_h (e^{\delta T} - 1)) + \beta T - \delta T \right) \\ &\quad + \frac{1}{\psi_l} \log \left( \frac{L \psi_h \psi_l (e^{\beta T} - 1)}{\beta (\psi_h + \psi_l)} \right) - T \beta \end{aligned} \quad (35)$$

$$\begin{aligned} x_h^* &= \frac{\alpha_h}{\psi_h} \left( -\log(\gamma \delta L \psi_h \psi_l (e^{\beta T} - 1)) + \log(\beta \psi_l (e^{\delta T} - 1)) + \beta T - \delta T \right) \\ &\quad + \frac{1}{\psi_h} \log \left( \frac{L \psi_h \psi_l (e^{\beta T} - 1)}{\beta (\psi_h + \psi_l)} \right) - T \beta \end{aligned} \quad (36)$$

$$\eta_l^* = \left( \frac{\beta (e^{\delta T} - 1) e^{\alpha_l (\log(\gamma \delta L \psi_h \psi_l (e^{\beta T} - 1)) - \log(\beta \psi_h (e^{\delta T} - 1)) + \beta(-T) + \delta T) + T(\beta - \delta)}}{\gamma \delta L \psi_l (e^{\beta T} - 1)} \right)^{\frac{1}{1-\alpha_l}} \quad (37)$$

$$\eta_h^* = \left( \frac{\beta (e^{\delta T} - 1) e^{T(\beta - \delta) + \alpha_h (\log(\gamma \delta L \psi_h \psi_l (e^{\beta T} - 1)) - \log(\beta \psi_l (e^{\delta T} - 1)) - \beta T + \delta T)}}{\gamma \delta L \psi_h (e^{\beta T} - 1)} \right)^{\frac{1}{1-\alpha_h}} \quad (38)$$

where  $z^* = z^\diamond$ . Assuming that  $\alpha_{i \in \{l, h\}} > 0$ ,  $\psi_{i \in \{l, h\}} > 0$ ,  $L > 0$ ,  $T > 0$ ,  $\gamma > 0$ ,  $\delta > 0$  and  $\beta > 0$ , then Equations 35 and 36 simplify to the result given in Proposition 1 (Part 1) and Equations 37 and 38 simplify to the equations given in Proposition 1 (Part 2).  $\square$

## A.2 Proof of Proposition 2

**Part 1: Target Investment with Steward** For the fully informed steward, setting  $\bar{x}_{i \in \{h, l\}}$  and  $\bar{z}$  the steward's objective is to minimize total aggregate loss for all targets. For our derivation, the targets are all assumed to be identical therefore the steward seeks to minimize

$$\tilde{V}_P = N_T \int_{t_0}^T e^{-\bar{\beta} t} (z \tilde{\sigma}(x_l, \eta_l^\diamond) + (1-z) \tilde{\sigma}(x_h, \eta_h^\diamond)) dt + N_T x_h + N_T x_l$$

where  $\tilde{\sigma}_{i \in \{l, h\}} : \mathbb{R}_+ \rightarrow [0, 1]$ ,  $\eta^\diamond_{i \in \{l, h\}}$  is derived from Equations 33 and 34. The asset allocation  $z$  does not have a tractable analytic solution in this case, so for exposition purposes we focus on  $x_l$  and  $x_h$  when  $z$  is fixed. In this case, let us fix  $z$  to the Nash equilibrium solution, therefore  $\bar{z} = z^\diamond$ , from the proof in Proposition 1 (Part 1). Evaluating the integral from  $t_0 = 0$  to  $T$  and eliminating  $N_T$  yields:

$$\begin{aligned} \tilde{V}_P &= \frac{L(e^{\beta T} - 1) e^{-x_h \psi_h - x_l \psi_l - \beta T} \psi_h e^{x_l \psi_l}}{\beta(\psi_h + \psi_l)} \left( \frac{\psi_h (e^{\delta T} - 1) e^{-x_h \psi_h - T \delta}}{\gamma \delta (\psi_h + \psi_l)} \right)^{\frac{\alpha_h}{1 - \alpha_h}} \\ &+ \frac{L(e^{\beta T} - 1) e^{-x_h \psi_h - x_l \psi_l - \beta T} \psi_h e^{x_l \psi_l} \psi_l e^{x_h \psi_h}}{\beta(\psi_h + \psi_l)} \left( \frac{\psi_l (e^{\delta T} - 1) e^{-x_l \psi_l - \delta T}}{\gamma \delta (\psi_h + \psi_l)} \right)^{\frac{\alpha_l}{1 - \alpha_l}} \end{aligned} \quad (39)$$

This is now a two-dimensional unconstrained optimization problem, where

$$\frac{\partial \tilde{V}_P}{\partial x_l} = \frac{L \psi_l^2 (e^{\beta T} - 1) e^{-x_l \psi_l - \beta T}}{\beta(\alpha_l - 1)(\psi_h + \psi_l)} \left( \frac{\psi_l (e^{\delta T} - 1) e^{-x_l \psi_l - \delta T}}{\gamma \delta (\psi_h + \psi_l)} \right)^{\frac{\alpha_l}{1 - \alpha_l}} \quad (40)$$

$$\frac{\partial \tilde{V}_P}{\partial x_h} = \frac{L \psi_h^2 (e^{\beta T} - 1) e^{-x_h \psi_h - \beta T}}{\beta(\alpha_h - 1)(\psi_h + \psi_l)} \left( \frac{\psi_h (e^{\delta T} - 1) e^{-x_h \psi_h - \delta T}}{\gamma \delta (\psi_h + \psi_l)} \right)^{\frac{\alpha_h}{1 - \alpha_h}}. \quad (41)$$

Setting  $\partial \tilde{V}_P / \partial x_l = 0$  and  $\partial \tilde{V}_P / \partial x_h = 0$  and solving for  $\bar{x}_l$  and  $\bar{x}_h$ , we obtain the steward's solution:

$$\bar{x}_l = \frac{-(1 - \alpha_l)}{\psi_l} \times \quad (42)$$

$$\log \left( \frac{(1 - \alpha_l) \beta \gamma^{-\frac{\alpha_l}{\alpha_l - 1}} \delta^{-\frac{\alpha_l}{\alpha_l - 1}} \psi_l^{\frac{1}{\alpha_l - 1} - 1} (\psi_h + \psi_l)^{\frac{1}{1 - \alpha_l}} (e^{\delta T} - 1)^{\frac{1}{\alpha_l - 1} + 1} e^{\beta T - \frac{\delta \alpha_l}{\alpha_l - 1} T}}{L(e^{\beta T} - 1)} \right)$$

$$\bar{x}_h = \frac{-(1 - \alpha_h)}{\psi_h} \times \quad (43)$$

$$\log \left( \frac{(1 - \alpha_h) \beta \gamma^{-\frac{\alpha_h}{\alpha_h - 1}} \delta^{-\frac{\alpha_h}{\alpha_h - 1}} \psi_h^{\frac{1}{\alpha_h - 1} - 1} (\psi_h + \psi_l)^{\frac{1}{1 - \alpha_h}} (e^{\delta T} - 1)^{\frac{1}{\alpha_h - 1} + 1} e^{T(\beta - \frac{\delta \alpha_h}{\alpha_h - 1})}}{L(e^{\beta T} - 1)} \right)$$

Simplification of Equations 42 and 43 yields the solutions given in Proposition 2 (Part 1).

**Part 2: Attacking Intensity** The attacker intensities under the fully informed steward are obtained by substituting the optimal expenditures  $\bar{x}_l$  and  $\bar{x}_h$  into Equations 33 and 34; that is,

$$\bar{\eta}_l = \left( \frac{\psi_l (e^{\delta T} - 1) e^{\delta(-T) - \bar{x}_l \psi_l}}{\gamma \delta (\psi_h + \psi_l)} \right)^{\frac{1}{1 - \alpha_l}} \quad (44)$$

$$\bar{\eta}_h = \left( \frac{\psi_h (e^{\delta T} - 1) e^{\delta(-T) - \bar{x}_h \psi_h}}{\gamma \delta (\psi_h + \psi_l)} \right)^{\frac{1}{1-\alpha_h}} \quad (45)$$

Setting  $i = \{h, l\}$  and  $j = \{h, l\}$  for  $j \neq i$  yields Equation 9 in Proposition 2 (Part 2).  $\square$

The analytic forms of Equations 44 and 45, as functions of the model parameters, are as follows:

$$\bar{\eta}_l = \left( \frac{\psi_l e^{\delta(-T)} (e^{\delta T} - 1)}{\gamma \delta (\psi_h + \psi_l)} \right)^{\frac{1}{1-\alpha_l}} \times \quad (46)$$

$$\left( - \frac{\beta (\alpha_l - 1) \gamma^{-\frac{\alpha_l}{\alpha_l-1}} \delta^{-\frac{\alpha_l}{\alpha_l-1}} \psi_l^{\frac{1}{\alpha_l-1}-1} (\psi_h + \psi_l)^{\frac{1}{1-\alpha_l}} (e^{\delta T} - 1)^{\frac{1}{\alpha_l-1}+1} e^{T(\beta - \frac{\delta \alpha_l}{\alpha_l-1})}}{L (e^{\beta T} - 1)} \right)$$

$$\bar{\eta}_h = \left( \frac{\psi_h e^{\delta(-T)} (e^{\delta T} - 1)}{\gamma \delta (\psi_h + \psi_l)} \right)^{\frac{1}{1-\alpha_h}} \times \quad (47)$$

$$\left( - \frac{\beta (\alpha_h - 1) \gamma^{-\frac{\alpha_h}{\alpha_h-1}} \delta^{-\frac{\alpha_h}{\alpha_h-1}} \psi_h^{\frac{1}{\alpha_h-1}-1} (\psi_h + \psi_l)^{\frac{1}{1-\alpha_h}} (e^{\delta T} - 1)^{\frac{1}{\alpha_h-1}+1} e^{T(\beta - \frac{\delta \alpha_h}{\alpha_h-1})}}{L (e^{\beta T} - 1)} \right).$$

### A.3 Proof of Proposition 4

The final case we consider in this paper considers the case when a steward can only observe and mandate one of the elements of the investment allocation,  $x_h$ . The targets have discretion to signal a value  $\hat{L}$ , however the steward does not know the true value of  $L$  or  $z$ .

The attackers signal a value  $\tilde{\zeta}$ , which we assume is actually  $1 - z$ . Targets, still have to choose their asset allocation, but they can potentially hide it from a potentially costly investment allocation. For tractability, we will assume this is in two steps, a signal of  $\hat{L}$  and  $\tilde{\zeta}$  and then an adjustment. This is done for tractable exposition, although the simultaneous model also has an analytic solution and provides a similar result, whilst being algebraically more complex.

**Part 1: Asset Class  $h$**  Let  $\tilde{\sigma}_{i \in \{l, h\}} : \mathbb{R}_+ \rightarrow [0, 1]$  and, for the targets, let  $x_h$  be exogenous. Targets minimize

$$\tilde{V}_T = \frac{1}{\beta} L (e^{\beta T} - 1) e^{-x_h \psi_h - x_l \psi_l + \beta(-T)} (z e^{x_h \psi_h} \eta_l^{\alpha_l} - (z - 1) \eta_h^{\alpha_h} e^{x_l \psi_l}) + x_h + x_l \quad (48)$$

Setting  $\partial \tilde{V}_T / \partial x_l = 0$  and  $\partial \tilde{V}_T / \partial z = 0$  and solving for  $x_l$  and  $z$  we obtain

$$x_l^\diamond = \frac{x_h \psi_h - \log(\eta_h^{\alpha_h} \eta_l^{-\alpha_l})}{\psi_l} \quad (49)$$

$$z^\diamond = \frac{\beta e^{\beta T} \eta_l^{-\alpha_l}}{L \psi_l (\eta_h^{\alpha_h} \eta_l^{-\alpha_l} e^{\beta T - x_h \psi_h} - \eta_h^{\alpha_h} e^{-x_h \psi_h} \eta_l^{-\alpha_l})} \quad (50)$$

Note that the both the optimal asset allocation  $z^\diamond$  and the optimal investment  $x_i^\diamond$  are now functions of  $x_h$  and are both subject to an upper bound of  $\eta_i^* < e^{\alpha_i^{-1} x_i \psi_i}$ .

**Part 2: Asset Class  $l$**  The steward has received a information on  $\hat{L}$  and  $\tilde{\zeta}$ , which in this derivation we treat as exogenous. However, the optimal initial bid of  $\hat{L}$  from the targets to the steward can be obtained by numerical analysis. The steward sets a mandatory investment level of  $\bar{x}_h$ , from a restricted information set by minimizing

$$\tilde{V}_P = \frac{N_T}{\bar{\beta}} \hat{L} \left( e^{\bar{\beta}T} - 1 \right) \eta_h^{\diamond \alpha_h} e^{-\bar{\beta}T - x_h \psi_h} + N_T x_h \quad (51)$$

where

$$\eta_h^\diamond = \left( \frac{\tilde{\zeta} (e^{\delta T} - 1) e^{-x_h \psi_h - \delta T}}{\gamma \delta} \right)^{\frac{1}{1 - \alpha_h}} \quad (52)$$

solving the single equation and single unknown  $\partial V_P / x_h = 0$ , yields

$$\bar{x}_h = \frac{1}{\psi_h} \log(\mathcal{A}) + \frac{\alpha_h}{\psi_h} (T(\bar{\beta} - \delta) - \log(\mathcal{B})) - \frac{\bar{\beta}T}{\psi_h} \quad (53)$$

Note that  $x_h$  is now a function of  $\hat{L}$ ,  $\tilde{\zeta}$  and the structural parameters  $\delta$ ,  $\gamma$ ,  $\psi_{i \in \{l, h\}}$ ,  $\alpha_{i \in \{l, h\}}$  and  $T$ . Simplification of Equation 53 results in the steward component of Proposition 4 (Part 1). Substitution of  $\bar{x}_h$  into Equation 52 provides the functional form of the attacker intensity  $\bar{\eta}_h$  of Proposition 4 (Part 1). The solution in terms of the model parameters is as follows:

$$\bar{\eta}_h = \mathcal{B}^{\frac{1}{1 - \alpha_h}} e^{\alpha_h (\log(-\mathcal{A}) + T(\delta - \bar{\beta}) + T(\bar{\beta} - \delta))} \frac{1}{1 - \alpha_h} \quad (54)$$

where

$$\mathcal{A} = \frac{\hat{L} (1 - e^{bT}) \psi_h \gamma^{\frac{1}{\alpha_h - 1} + 1} \delta^{\frac{1}{\alpha_h - 1} + 1} \tilde{\zeta}^{\frac{\alpha_h}{1 - \alpha_h}} (e^{\delta T} - 1)^{\frac{\alpha_h}{1 - \alpha_h}}}{\bar{\beta} (\alpha_h - 1)} \quad (55)$$

$$\mathcal{B} = \frac{\bar{\beta} (1 - \alpha_h) \gamma^{\frac{1}{1 - \alpha_h} - 2} \delta^{\frac{1}{1 - \alpha_h} - 2} \tilde{\zeta}^{\frac{1}{\alpha_h - 1} + 2} (e^{\delta T} - 1)^{\frac{1}{\alpha_h - 1} + 2}}{\hat{L} (e^{bT} - 1) \psi_h} \quad (56)$$

To derive the target allocation and attacker intensity  $\eta_l^\ddagger$ , we now simply need to substitute the functional forms of  $\bar{x}_h$  and  $\bar{\eta}_h$  into Equations 49 and 50 and simplify functional forms in Proposition 4 (Part 2). For  $x_l^\ddagger$ ,  $z^\ddagger$  and  $\eta_l^\ddagger$ ,

$$x_l^\ddagger = \frac{1}{\psi_l} \alpha_h (T(\bar{\beta} - \delta) - \log(\mathcal{A})) + \log(\mathcal{A}) - \bar{\beta}T \quad (57)$$

$$- \frac{1}{\psi_l} \log \left( \left( \frac{\beta (e^{\delta T} - 1) e^{T(\beta - \delta)}}{\gamma \delta L \psi_l (e^{\beta T} - 1)} \right)^{-\alpha_l} \left( e^{\alpha_h (\log(\mathcal{A}) - \bar{\beta}T + \delta T) + T(\bar{\beta} - \delta)} \right)^{\frac{\alpha_h}{1 - \alpha_h}} \right)$$

$$z_l^\ddagger = \frac{\beta \mathcal{A}}{L \psi_l (e^{\beta T} - 1)} e^{T(\beta - \bar{\beta}) - \alpha_h (\log(\mathcal{A}) + T(\delta - \bar{\beta}))} \times \quad (58)$$

$$\left( \mathcal{B} e^{\alpha_h (\log(\mathcal{A}) + T(\delta - \bar{\beta})) + T(\bar{\beta} - \delta)} \right)^{\frac{-\alpha_h}{1 - \alpha_h}}$$

$$\eta_l^\ddagger = \frac{\beta (e^{\delta T} - 1) e^{T(\beta - \delta)}}{\gamma \delta L \psi_l (e^{\beta T} - 1)} \quad (59)$$

which simplify to the equations in Proposition 4 (Part 2).  $\square$