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Abstract

We present a general framework to tackle the problem of finding time-independent dynamics generating target unitary evolutions. We show that this problem is equivalently stated as a set of conditions over the spectrum of the time-independent gate generator, thus translating the task into an inverse eigenvalue problem. We illustrate our methodology by identifying suitable time-independent generators implementing Toffoli and Fredkin gates without the need for ancillae or effective evolutions. We show how the same conditions can be used to solve the problem numerically, via supervised learning techniques. In turn, this allows us to solve problems that are not amenable, in general, to direct analytical solution, providing at the same time a high degree of flexibility over the types of gate-design problems that can be approached. As a significant example, we find generators for the Toffoli gate using only diagonal pairwise interactions, which are easier to implement in some experimental architectures. To showcase the flexibility of the supervised learning approach, we give an example of a non-trivial four-qubit gate that is implementable using only diagonal, pairwise interactions.

1. Introduction

Let us consider the synthesis of a quantum operation $G$ (a gate) from the underlying dynamics of a quantum system. Unitarity of $G$ implies the existence of a Hermitian generator $H_G$ such that $G = e^{i H_G}$ (we assume units such that the simulation time $t$ is dimensionless, and rescale it so that the desired gate is successfully achieved at $t = 1$). Such generator will typically contain highly non-local interactions, which can be difficult to realize in a given physical setup. However, for a choice of the platform to use for the implementation of $G$, it is generally possible to single out a subset $\Gamma$ of physical interactions that can be realized relatively easily and inexpensively. The question that we address here is thus: is it possible to synthesize $G$ from an $H_G$ comprising operations drawn only from an assigned $\Gamma$?

This question is very relevant, for instance, in the context of quantum simulation, where the problem of general reachability of a target dynamics given a set of allowed physical interactions is pivotal [1]. However, it is also important for the realization of large-scale quantum computation [2, 3], which relies on the capability of implementing entangling gates between many qubits with high fidelity. A notable case is the quantum Toffoli gate, which is the quantum counterpart of the universal reversible classical Toffoli gate [4]. When paired with a single-qubit gate, the quantum Toffoli gate is universal for quantum computation [5].
Moreover, it is optimal for quantum error correction [6–9], and is a key component for reversible arithmetic operations such as modular exponentiation [10]. Unfortunately, the natural dynamics generating Hilbert spaces to overcome these limitations include a suitable use of the additional processing power offered by larger Hilbert spaces encompassing ancillary information carriers [11], and the embedding of quantum control techniques [12, 13]. The identification of suitable alternatives to such expensive strategies for gate synthesis and simulation would represent a significant contribution to the ongoing effort towards the translation of theoretical protocols to the production line of quantum technologies [14]. It is worth stressing that the strategies put forward in this paper differ substantially from techniques such as quantum control and gate compilation. Quantum control techniques allow for time-dependent Hamiltonians, whereas gate compilation foregoes dealing directly with Hamiltonians, looking instead for sequences of gates whose combined action results in the target operation. We, on the other hand, aim to find strategies that involve a single, time-independent Hamiltonian. These are sometimes referred to as single-shot strategies. This constraint changes the task substantially, making most techniques used in the context of quantum control not seamlessly applicable. On the other hand, gate compilation requires the underlying assumption that the elements of the set to be used in order to decompose a given unitary are all reliably implementable in the architecture under consideration. This fully bypasses any concern about the dynamics necessary to actually implement a given gate. In addition, gate compilation is combinatorial in nature. Our approach to the learning of a quantum gate, instead, makes use of a set of continuous parameters which determine the final operation only through a matrix exponential operation.

In this paper we show how the interplay between analytical results and efficient numerics—as enabled by the supervised learning paradigm—provides a powerful and flexible tool to explore a wide variety of gate design and simulation scenarios.

More specifically, we identify three analytical conditions that, when met, provide a Hamiltonian \( \tilde{H} \) comprising only operations drawn from \( \Gamma \) and such that \( G = \exp(i\tilde{H}) \) for a given target \( G \). While these conditions do in principle allow to solve the problem in its generality, the corresponding equations soon become practically intractable. Even in this case, however, our framework provides an improved ansatz, which can be used to perform numerical optimization via supervised learning.

For the numerical optimization, we showcase an original application of supervised machine learning that allows to obtain results more efficiently than alternative methods. The context in which supervised learning is commonly explored is roughly the following: given a training dataset consisting of pairs \((x_i,\ell_i)\) in \( X \times L \) and a parametrized function (usually referred to as the model) \( f_\lambda : X \mapsto L \), find \( \lambda_0 \) such that \( f_{\lambda_0}(x_i) \approx \ell_i, \forall i \). This is usually done by minimizing the empirical loss function, that is, by solving the optimization problem \( \arg\min_\lambda \{ \sum_i L(f_\lambda(x_i), \ell_i) \} \) for some loss function \( L \) quantifying the distance between expected and obtained outputs (typical choices being the \( L_2 \) distance and the delta function). In other words, the goal is to approximate an unknown \( g : X \mapsto L \) from a finite number of samples. In this sense, supervised learning is akin to function fitting. Notably, this is not the kind of problem naturally encountered in gate design. The design of a time-independent Hamiltonian generating a target \( G \), given a parametrization \( H_\lambda \) for the Hamiltonian, fits into the standard scenario of solving argmin\( \lambda \{ L(\lambda) \} \) with \( L(\lambda) \) the operational distance between \( \exp(iH_\lambda) \) and \( G \). However, as we will show here, supervised learning techniques can be fruitfully applied to this kind of problems.

Optimizing \( L(\lambda) \) is a daunting task, as the evaluation of \( e^{i\tilde{H}} \) requires up to \( O(d^4) \) operations [15] for each iteration (\( d \) is the dimension of the Hilbert space). Nonetheless, one finds that \( L(\lambda) \) can be expressed as an average over random states [16]. This suggests the use of stochastic optimization techniques, commonly employed in machine learning, where one optimizes an estimated empirical loss function, rather than the exact one. For gate design, the resulting estimated loss is \( L(\lambda) \approx \sum D \left( \exp \left( iH_\lambda \right) \left| \psi_i \right\rangle \left\langle \psi_i \right| G \right) \) with \( D \) a state distance and \( \left| \psi_i \right\rangle \) a set of random states. As in supervised learning, the latter are used to train the quantum system so that its dynamics reproduces the expected action \( G \). As typically \( H_\lambda \) is sparse, the computation of \( e^{i\tilde{H}} \) is efficient [17], unlike the full operator exponential. Moreover, stochasticity also helps to avoid local minima. When an analytic solution is not available, our framework gives a lower-dimensional representation of the loss function, which is then optimized via stochastic gradient descent (SGD) [18, 19]. Such representation is then embedded into an automatic differentiation (AD) protocol [20–23], to efficiently compute the gradients without numerical approximations. This significantly increases the efficiency of the optimization, allowing for the exploration of previously out-of-reach scenarios, and a systematic analysis of gate design and simulation problems. While the use of SGD to tackle gate design was also present in [16], our work achieves both analytical and numerical steps forwards. In particular, no analytical framework was developed in [16], and thus exact solutions could not
We start by computing an Hamiltonian $\hat{H}_G$ generating the target gate $G = e^{iHG}$. Using the spectral decomposition of $G$, we have $\hat{H}_G = -iU \log(\Lambda)U^\dagger$, where $\hat{G} = UGU^\dagger$, $\log$ denotes the principal branch of the logarithm, and $\Lambda$ is the diagonal matrix with the eigenvalues of $U$. Fixing a branch for the logarithm makes it single-valued, and $\hat{H}_G$ uniquely determined from $\hat{G}$. To distinguish this particular Hamiltonian from the other possible generators for $G$, we will refer to it as the principal generator. In general, the generator $\hat{H}_G$ will contain both physical interactions, that can be realized easily in a given physical setup, and unphysical ones, i.e. dynamics that are not naturally achieved in the chosen platform. Our goal is to construct a new Hamiltonian $\tilde{H}_G$, comprising only physical interactions, such that $G = \exp(i\tilde{H}_G) = \exp(i\hat{H}_G)$.

We assume that $\tilde{H}_G$ depends on a quantity $\lambda$ that parametrizes the set of physical interactions to be used. For instance, in a spin system, the physical interactions can be a certain subset of the possible two-body and single-body interactions, like the set of coupling strengths and local magnetic fields. In general, $\tilde{H}_G(\lambda)$ may also model a system where the register is coupled to auxiliary degrees of freedom, though we will focus on the case without ancillary qubits. The following three conditions are necessary and sufficient for $\tilde{H}_G$ to satisfy our requirements

$$\tilde{H}_G \text{ contains only physical interactions,}$$

$$[\tilde{H}_G, \hat{H}_G] = 0,$$

$$\text{Eig}(\tilde{H}_G - \hat{H}_G) = \{2\pi n_i\} \quad (n_i \in \mathbb{Z}).$$
Equations (1b) and (1c) ensure that $G = \exp(i\hat{H}_G - i\hat{H}_G) \exp(i\hat{H}_G)$ and $\exp(i\hat{H}_G - i\hat{H}_G) = 1$, while (1a) reiterates the constraints on $\hat{H}_G$. While equation (1b) may seem too restrictive, this turns out to not be the case. To see this, consider the decomposition $G = \sum \lambda_k \sum_j P_{kj}$ of a general gate $G$. Here, $P_{kj}$ is the $j$th trace-1 projector in the $k$th degenerate subsector of the eigenspace. It follows that $\hat{H}_G$ and $\hat{H}_G$ can be written as $\hat{H}_G = -i\sum_k \log(\lambda_k) I_k$ and $\hat{G} = H_G + 2\pi i\sum_{k, l} \nu_{kl} P_{kl}$. Here $I_k = \sum_j P_{kj}$, which is not (in general) an identity (or diagonal) matrix, and we have used the expression $\log \lambda = \log \lambda + 2\pi i\nu$ ($\nu \in \mathbb{Z}$), applied to every term of the spectral decomposition of $G$. For any $\hat{H}_G$ we thus have $[\hat{H}_G, \hat{G}] = 0$.

Equations (1a–1c) considerably simplify the problem of gate synthesis, and can be used in several ways. First, they can be analytically solved in at least some situations of physical interest. Second, they can be used to produce an efficient starting point for numerical optimization techniques. The general procedure is to start from a parametrized expression for $\hat{H}_G(\lambda)$ satisfying the condition stated in equation (1a), and then use equation (1b) to both reduce the set of possible interactions and impose constraints on the parameters. The problem then reduces to the enforcing of the constraints set by equation (1c). This is the non-trivial step in the procedure, which we will however show to be analytically solvable in at least some cases. This strategy thus reduces the task of constrained gate design into an inverse eigenvalue problem, a topic well studied in the field of numerical analysis [28]. More generally, we develop a numerical supervised learning technique to avoid having to solve directly the non-trivial eigenvalue problem posed by equation (1c). It is worth noting that, while $\hat{H}_G$ produces the same unitary evolution given by $G$ at $t = 1$, the dynamics will in general be different at $0 < t < 1$.

3. Applications: Toffoli and Fredkin gates

3.1. Quantum Toffoli gate

The quantum Toffoli gate $U_{\text{Toff}}$ is a control-control-NOT that flips the state of the target qubit (the third qubit in our case) when the state of the two controls (first and second qubits) is $|1\rangle_1 \otimes |1\rangle_2$, and acts trivially on the target qubit otherwise. Its realization is an important step towards the construction of quantum computers [9, 29–31]. A time-independent two-body Hamiltonian that simulates $U_{\text{Toff}}$ with four qubits has been obtained in [16] using numerical optimization, while three qubits have only been found to make approximate and classical Toffoli gates [32]. Here, following the construction in equation (1), we find an analytic solution that requires only three qubits.

The principal generator of $U_{\text{Toff}}$, obtained by taking the principal value of its logarithm, is

$$\hat{H}_{\text{Toff}} = (\pi/8)(1 - \sigma_j^z)(1 - \sigma_j^z)(1 - \sigma_j^z),$$

where only three-qubit term is $\sigma_j^z\sigma_j^z\sigma_j^z$, where $\sigma_j^\alpha$ is the $\alpha$ Pauli matrix acting on the $j$th qubit. We now parametrize $\hat{H}_{\text{Toff}}$ as

$$\hat{H}_{\text{Toff}} = h_0 \mathbb{1} + \sum_i h_i^\alpha \sigma_i^\alpha + \sum_{i,j} f_{ij}^{\alpha\beta} \sigma_i^\alpha \sigma_j^\beta,$$

where $h_0, h_i^\alpha, f_{ij}^{\alpha\beta}$ are real parameters. The above expression, containing 37 parameters, automatically satisfies equation (1a) in that it corresponds to an $\hat{H}_{\text{Toff}}$ without three-qubit interactions. Imposing the condition in equation (1b) removes 12 parameters, leaving us with 25 of them. These are still too many to solve the inverse eigenvalue problem of equation (1c). We thus impose physically plausible assumptions on the coefficients in order to obtain a generator with a small enough number of parameters, so that equation (1c) is satisfied and the resulting equations solved. In particular, we set

$$f_{11}^{\alpha\beta} = f_{12}^{\alpha\beta} = f_{13}^{\alpha\beta} = f_{22}^{\alpha\beta} = 0, \quad f_{13}^{\alpha\beta} = f_{23}^{\alpha\beta} = \pi/8, \quad f_{23}^{\alpha\beta} = -f_{13}^{\alpha\beta}, \quad h_1^\alpha = h_2^\alpha = -\pi/8.$$

The rationale behind these assumptions is to look for a generator that is diagonal with respect to the first two qubits, does not use $\sigma_j^\alpha$ operators, and does not introduce new off-diagonal interactions, on top of the ones already in the principal generator. This last assumption is useful because it implies a reduced number of parameters in $\hat{H}_{\text{Toff}} = \hat{H}_{\text{Toff}} - \hat{H}_{\text{Toff}}$, which is the operator on which we have to impose equation (1c).

The projectors are such that $\sum_{k} P_{kl} = \mathbb{1}_4, P_{0j} P_{0j} = 0$ if $j \neq f$, and $P_{j} P_{k} = 0$ if $k \neq k'$ for any $j, f$. 

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Note that different assumptions, leading to different solutions, are possible. Imposing the above constraints we obtain

$$\mathcal{H}^{'}_{\text{Toff}} = (\pi/8)\sigma_1^z\sigma_2^z\sigma_3^z + (h_0 - \pi/8)\mathbb{I} + (h_3^z + \pi/8)\sigma_3^z + (J_{13}^z - \pi/8)\sigma_1^z\sigma_2^z + J_{13}^y(\sigma_1^z - \sigma_2^z)\sigma_3^z. \quad (5)$$

The problem is now to find values for the coefficients in equation (5) such that \(\exp(i\mathcal{H}^{'}_{\text{Toff}}) = 1\), which is equivalent to finding coefficients such that all the eigenvalues of \(\mathcal{H}^{'}_{\text{Toff}}\) are integer multiples of 2\(\pi\). Solving for \(h_0, h_3^z, J_{13}^z, J_{13}^y\) gives several solutions. For example, we find that

$$\tilde{\mathcal{H}}_{\text{Toff}}(\nu) = \frac{\pi}{8} \left[ 1 + 4|\nu| - 2\sigma_3^z - \sigma_1^z - \sigma_2^z + (\sigma_1^z + \sigma_2^z)\sigma_3^z + (1 - 4|\nu|)\sigma_1^z\sigma_2^z + \sqrt{16\nu^2 - 1}(\sigma_2^z - \sigma_1^z)\sigma_3^z \right]$$

is a class of generators for \(\mathcal{U}_{\text{Toff}}\), satisfying \(\exp(i\tilde{\mathcal{H}}_{\text{Toff}}(\nu)) = \mathcal{U}_{\text{Toff}}\) for all non-zero integers \(\nu\). In appendix A3 we present a larger class of generators for the Toffoli, as well as provide an in-depth analysis of how these generators produce the Toffoli gate. Moreover, in appendix A4, we discuss a possible line of reasoning to arrive to one of these generators via direct algebraic manipulation. While with this ad hoc reasoning we cannot recover the full set of generators obtainable via our framework, we still gather insight on the way non-trivial interaction coefficients emerge. Moreover, it showcases the type of sophisticated analysis needed to solve these problems without using our framework. Thus, a highly non-trivial gate such as Toffoli’s, which in principle requires three-body interactions as in equation (2), can be obtained exactly without three-qubit interactions. Notably, although the generators for the Toffoli gate found here contain non-diagonal \(\sigma^z\)–\(\sigma^x\) interactions, which may be hard to implement, we will show in section 4 that this can be overcome using a supervised learning approach.

3.2. Quantum Fredkin gate

On a similar note, it is possible to use the framework provided by equations (1a–1c) to find a Hamiltonian that does not contain three-qubit interaction terms, and generates the Fredkin gate at suitable times. The quantum Fredkin gate \(\mathcal{U}_{\text{Fred}}\) is a three-qubit gate which swaps two qubits conditionally to the first qubit being in the \(|1\rangle\) state, and is of use for a number of quantum information protocols [33, 34]. A time-independent two-body Hamiltonian simulating \(\mathcal{U}_{\text{Fred}}\) with four qubits has been found in [16] using numerical optimization. We find an analytic solution that requires as few as three qubits. Explicitly

$$\mathcal{H}_{\text{Fred}} = \frac{\pi}{8} \left( 1 - \sigma_1^z \right) \left[ 1 - \sum_{\alpha} \sigma_\alpha^z \sigma_\alpha^z \right], \quad (7)$$

where qubit 1 is the control. We now write down the general parametrization \(\tilde{\mathcal{H}}_{\mathcal{G}}(\lambda)\) for a three-qubit Hamiltonian containing only pairwise diagonal interactions, and imposing equation (1b) we cut the number of parameters \(\lambda\) down to 22. Imposing additional physically plausible conditions, like the symmetry of second and third qubit, we manage to reduce the number of parameters enough to solve the eigenvalue problem, finding

$$\tilde{\mathcal{H}}_{\text{Fred}} = \frac{\pi}{8} \left( \sqrt{\frac{143}{5}} \mathbb{I} + 5\sqrt{3}\sigma_1^z \right) (\sigma_2^z + \sigma_3^z) - \frac{3\pi}{8} \sum_{\alpha=x,y,z} \sigma_\alpha^z \sigma_\alpha^z - \frac{3\pi}{4} \sqrt{\frac{7}{5}} \sigma_1^z (\sigma_2^z + \sigma_3^z) + \frac{\pi}{2} \sigma_1^z + \frac{3\pi}{8} \mathbb{1}. \quad (8)$$

Therefore, also a non-trivial gate of crucial relevance such as the Fredkin gate can be implemented without time-dependent dynamics using only at most two-qubit interactions. The physical reason behind this simplification can be understood from the study of the spectral properties. For example, the gate \(\mathcal{U}_{\text{Fred}}\) has only two eigenvalues, \(\lambda_\pm = \pm 1\) with \(\lambda_+\) having a sevenfold degeneracy. Such degeneracy makes the propagator generated by \(\mathcal{H}_{\text{Fred}}\) operate in a two-level subspace. On the other hand, the spectrum of \(\mathcal{H}_{\text{Fred}} - \mathcal{H}_{\text{Fred}}\) is \([-4\pi, -2\pi, 0, 0, 0, 2\pi, 2\pi, 4\pi]\), showing that the degeneracy in the spectrum of \(\mathcal{H}_{\text{Fred}}\) is partially lifted when considering \(\mathcal{H}_{\text{Fred}} - \mathcal{H}_{\text{Fred}}\), and the dynamics thus occurs in an larger effective Hilbert space. Although \(\exp(i\mathcal{H}_{\text{Fred}})\) is non-symmetric for most of the evolution times, all the symmetries are restored at \(t = 1\) and any subsequent integer times. This shows that breaking the symmetries of \(\mathcal{G}\) and exploiting its degenerate space can help the gate design when restricting the set of viable interactions.
Algorithm 1. Stochastic gate design.

1: Choose values for the set $\lambda_0$ for $\mathcal{H}_0$, such that equation (1a) is satisfied.
2: Use equation (1b) to find a reduced set $\lambda$ (linearly related to $\lambda_0$).
3: repeat
   $\triangleright$ Iteratively solve equation (1c)
4: Generate a random set of $N_b$ input states $\{|\psi_i\rangle\}$ with $N_b$ the size of the mini-batches chosen beforehand.
5: For each $k$, compute $\nabla_{\lambda} F(\lambda, |\psi_i\rangle)$. Machine learning frameworks like Theano [38], TensorFlow [39], or PyTorch [40] (among others) enables the calculation of gradients automatically from the chain rule. This avoids numerical errors arising from numerical differentiation.
6: Update the coupling strengths $\lambda$. We do this using the so-called momentum gradient descent method [41], corresponding to the following updating rule
   $$v \rightarrow \gamma v + \eta \nabla_{\lambda} F(\lambda, |\psi_i\rangle), \quad \lambda \rightarrow \lambda + v.$$ (9)
   Here the learning rate $\gamma$ and the momentum term $\gamma$ are hyperparameters to be chosen beforehand. The value of $\eta$ should decrease with the iteration number.
7: until a satisfactory value of the fidelity is obtained.

4. Supervised learning approach

We now describe a different methodology to solve the difficult part of equation (1). While the direct algebraic approach fails as soon as we consider more than a few parameters, and for example already fails to find solutions for the Toffoli gate with only diagonal interactions, the method we present here scales better with the number of interactions and is easily generalized to any kind of structure of the qubit network. The idea is to adopt a supervised learning approach to solve the optimization problem of finding the set of Hamiltonian parameters generating a target evolution.

The problem we address can be stated as follows. Given a target gate $G$ and a parametrized Hamiltonian $\mathcal{H}(\lambda) = \sum_{i=1}^n \lambda_i O_i$, where $\lambda = \{\lambda_i\}$ is a set of real parameters and $O_i$ are Hermitian operators, we look for the set $\lambda_0$ such that $\exp(i\mathcal{H}(\lambda_0)) = G$. This can be reframed as an optimization problem by considering the fidelity function $F(\lambda, \psi_i) \equiv \langle \psi_i | G^\dagger \exp(i\mathcal{H}(\lambda)) | \psi_i \rangle$, for an arbitrary state $|\psi_i\rangle$. Clearly, $F(\lambda_0, \psi) = 1$ for all $|\psi\rangle$ iff $\exp(i\mathcal{H}(\lambda_0)) = G$. A possible approach to find such $\lambda_0$ is to consider the average fidelity function $F(\lambda)$, defined as the average of $F(\lambda, \psi)$ over all $\psi$. Explicit formulas for $F$ are known [35, 36], so that standard optimization methods can be used. This method is however inefficient for the problem at hand, due to the size of the parameter space. We thus turn to a different technique, exploiting how the fidelity landscape changes when changing $|\psi\rangle$ [16]. We use the fact that the only values of $\lambda$ for which the fidelity is 1 regardless of $|\psi\rangle$ are those corresponding to our solution. Using the SGD algorithm [18, 19, 37], we implement the procedure in algorithm 1. A more detailed presentation of this algorithm is given in appendix B. To find the interaction parameters implementing a Toffoli gate, using only one-qubit evolutions and two-qubit diagonal interactions (i.e. interactions of the form $\sigma_i^x \sigma_j^y$), we start the numerical training from the Hamiltonian obtained by imposing equation (1b) on the parametrized Hamiltonian containing the required interactions, which has the form

$$\mathcal{H}_{\text{Toff}} = \sum_{j=1,2} J_{jz}^2 (1 + \sigma_j^2) + J_{jz}^2 (\sigma_j^1 \sigma_j^2 + \sigma_j^2 \sigma_j^1) + J_{jz}^2 \sigma_j^1 \sigma_j^2$$

$$+ h_j^1 \sigma_j^1 + h_j^2 \sigma_j^2 + h_j^3 \sigma_j^3 + (J_{jz}^2 \sigma_j^1 + J_{jz}^2 \sigma_j^2)(1 + \sigma_j^3).$$ (10)

Starting from equation (10), many solutions can be found, depending on the chosen initial conditions and the random states that are used at each run. In figure 1(a), several solutions are shown, proving that it is possible to implement a Toffoli gate using only pairwise diagonal interactions. This analysis can be extended to a Fredkin gate, whose generator with only diagonal pairwise interactions and commuting with the principal generator of the Fredkin is found to have at most two-body terms. Using this model as starting point for the training we again obtain several solutions, some of which are shown in figure 1(b). We then test our supervised learning framework by exploring numerically the possibility of implementing Toffoli and Fredkin gates using a more restrictive set of interactions. The results provide a series of physically-feasible interaction parameters that realize Toffoli and Fredkin gates with good fidelities, as further detailed in appendix B4.

Note that in principle a more reliable figure of merit to evaluate the performance of a gate would be the diamond distance. However, in all the instances we considered, the two quantities are actually equivalent since we always obtain unit fidelity (within numerical errors).
Figure 1. Eight different sets of interaction parameters generating the Toffoli gate [panel (a)] and the Fredkin one [panel (b)]. For each shown solution the training was started from the ansatz provided by equation (10), and the analogous equation for the Fredkin, respectively. Panel (c): eight sets of interaction parameters for the ‘double Fredkin’ gate. Each one of the shown solutions corresponds to unit fidelity up to numerical precision (that is, fidelities greater than $1 \times 10^{-16}$). We refer to appendix B4 for the details of the optimizations. To use a more uniform the notation, we denote with $J_{ijk}$ the coefficient of the interaction parameter $\sigma_i^1 \sigma_j^2 \sigma_k^3$, and similarly for interactions between more than three qubits. For example, $J_{033} \equiv J_{zz2}$, $J_{200} \equiv \hbar$, and $J_{1010} \equiv J_{xx1}$.

There is no need to stick to the training of three-qubit networks. To illustrate this, we provide another example of successful application of our framework, this time to implement a non-trivial unitary evolution over four qubits. In particular, we train a four-qubit network to implement the doubly-controlled Fredkin gate $U_{FF} \equiv |0\rangle\langle0| \otimes U_{Fred} + |1\rangle\langle1| \otimes U_{Fred}^*$, where $U_{Fred}$ denotes a Fredkin gate where the control qubit is the third one, and the first two are the targets. Such a gate can be implemented using no more than two-qubit interactions, and such set can be further restricted to only diagonal ones. Some instances are shown in figure 1(c). Note that these results require no ad hoc approach or ansatz, differently from the approach used to derive equations (6) and (8). In particular, this simplifies the test of questions as ‘can gate $X$ be implemented using only a set of interaction $Y$’ without resorting to extra ad hoc calculations.

5. Conclusions

We have presented a general framework to approach constrained gate-synthesis problems. We have showed that the procedure is amenable to direct analytical solution, providing time-independent Hamiltonians generating Toffoli and Fredkin gates using only undemanding diagonal interactions and no ancillary qubits. To our knowledge, no previous attempt at such a decomposition has been reported so far. Generality can be added to our approach by powerful techniques of supervised learning of the interaction parameters, which allowed to find Hamiltonians with specified sets of interactions producing target unitary evolutions. Our approach and results are potentially of great interest to optimize experimental implementations of quantum algorithms in architectures such as linear optics and superconducting qubits. Moreover, it might open new avenues for the exploration of related problems such as Hamiltonian learning [42–44], which might be seen as addressing a question complementary to the one tackled here. In particular, it will be interesting to address whether it is possible to design a supervised-learning approach to the data available for a class of Hamiltonians that should be learnt, in line with the strategy put forward with our method, and thus devise a strategy that is competitive with respect to recently proposed Bayesian approaches [45].

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Appendix A. Analytical approach

In this section we provide a detailed description of how to apply the framework introduced in section 2 to approach gate design problems analytically.

In appendix A1 we show how our framework can be applied to tackle generic perfect state transfer problems. In appendix A2 we prove that a CNOT gate cannot be implemented using only single-qubit interactions. While this is an obvious result, the calculation can be interesting to show, in a simple case, how the eigenvalue conditions given in equation (1c) can be used to rule out that a gate can be implemented with a specific set of interactions. In section 3 we showed how to use our framework to derive Hamiltonians implementing the Toffoli gate using only one- and two-qubit interactions. In appendix A4 we will show how one of the solutions obtained in appendix A3 could have been obtained via direct analytical reasoning, without prior knowledge of our solution framework. This provides some insight into how the solutions actually generate the Toffoli gate, and illustrates the kind of ad hoc non-trivial reasoning that, without the use of equation (1), would have been necessary to find such solutions.

For notational convenience, we will in this section denote Pauli matrices with \( X_i, Y_i, Z_i \), instead of \( \sigma_x^i, \sigma_y^i, \) and \( \sigma_z^i \) as in the main text.

A1. Perfect state transfer

A one-dimensional quantum walk is described by the Hamiltonian \( H_W = \sum_{k=1}^{N-1} J_k |k\rangle\langle k-1| + B_k |k\rangle\langle k| \), where \( N \) is the length of the lattice upon which the walk takes place, \( J_k \) the transition rates between adjacent sites, \( B_k \) the local energies and \( |k\rangle \) defines the state where the ‘walker’ is at the \( k \)th site. A quantum walk Hamiltonian \( H_W \) admits perfect state transfer (PST) at time \( t \), i.e. the initial state of the walker initially at site \( k \) is perfectly retrieved at site \( N-j+1 \) if \( e^{-iH_W t} = \Xi \), where \( \Xi_{kj} = \delta_{k,N-j+1} \) is the reflection matrix. Necessary and sufficient conditions for PST are well understood \([24, 25]\): firstly, \( H_W \) has to be ‘mirror-symmetric’, that is \([H_W, \Xi] = 0\), and secondly the eigenvalues \( \{E_k\} \) of \( H_W \) should to satisfy the condition \( e^{\delta_{kj} t} = (\pm 1)^k \).

We show now that finding the parameters \( \{J_k, B_k\} \) for PST is a particular case of the Hamiltonian design problem for gate simulation. Specifically, the conditions for PST can be obtained from the construction in equation (1). In a 1D quantum walk, the ‘physical’ couplings are the nearest neighbour interactions, but \( H_\Xi = I \log(\Xi) = \pi(\Xi - 1)/2 \) is long range. Using \( H_W \) as \( H^{\text{phys}}(\lambda) \), where \( \lambda = \{J_k, B_k\} \), and defining \( H' = H_W - H_\Xi \) following condition equation (1c), where the physical interactions in \( H_\Xi \) have been reabsorbed into the definition of \( H_W \), one finds that: (i) condition (1b) is equivalent to the mirror-symmetry request \([H_W, \Xi] = 0\); (ii) from condition (1c) and the definitions of \( H_W \) and \( H_\Xi \), one finds that the spectrum of \( H_W \) satisfies \( e^{\delta_{kj} t} = (\pm 1)^k \)—indeed the eigenvalues of \( \Xi \) are \( \{0, \pi\} \), so \( E_k = \pi(2n_k + 1) \) for integer \( n_k \). It is straightforward to extend the above proof for more general transfers, such as with long-range interactions \([46]\), or for perfect fractionalal transfers \([47, 48]\).

A2. Proof that CNOT needs two-qubit terms

We here show how to use our framework to prove that a CNOT gate cannot be implemented using only one-qubit interactions. While this result is trivial, it is nonetheless interesting to show how the framework can be used to obtain this kind of impossibility results.

The spectral decomposition of the CNOT reads

\[
\text{CNOT} = Z_1^+ + Z_1^- X_2^+ - Z_1^- X_2^-, \quad (S1)
\]

where we made a canonical choice for the basis of the three-fold degenerate eigenspace corresponding to the eigenvalue +1, and defined \( Z_i^\pm \equiv (1 \pm Z_i)/2 \), and similarly for \( X_i^\pm \). More explicitly, we are considering the basis set \( \{Z_1^+, Z_1^-, Z_2^+, Z_2^-, Z_3^+, Z_3^-\} \) made out of trace-1 projectors for the degenerate space, whereas the fourth projector is bound to be \( Z_1^- X_2^- \). The corresponding principal Hamiltonian \( H_{\text{CNOT}} \), obtained by directly taking the logarithm of equation (S1), is

\[
H_{\text{CNOT}} = \pi Z_1^- X_2^- \cdot \quad (S2)
\]

Let us now consider what happens when the multivaluedness of the logarithm is taken into account, but no rotation of the degenerate eigenspace is performed. Considering only the factors with two-qubit interactions, the following expression is found:
\[
H_{\text{CNOT}}/2\pi \sim \nu_1 Z_1 Z_2 + 2\pi (1/2 + \nu_4) Z_1 X_2,
\]  

(S3)

where here \(\nu_i \equiv \nu_i - \nu_i\), and \(\nu_i \in \mathbb{Z}\) is the integer produced by application of the logarithm to the \(i\)th projector. Note how the \(Z_1 Z_2\) factor cannot be removed by any choice of \(\nu_i\), which could be interpreted as a proof that two-qubit interaction terms are indeed necessary to implement the CNOT gate. This, however, does not in principle preclude the possibility that an appropriate rotation of the degenerate space allows to obtain a generator with only local terms. To verify that this is not the case, we would have to consider a generic rotation \(R\) of the degenerate space, that is, an operator of the form \(R = \sum_{i=1}^3 r_i |i\rangle + 1_i\), with \(|+1_i\rangle\) the \(i\)th eigenvector in a fixed base of the degenerate space. The problem is thus to find a unitary \(R\) and integers \(\nu_i\) such that

\[
\nu_i R(Z_i^+ Z_j^+) R^\dagger + \nu_2 R(Z_i^+ Z_j^-) R^\dagger + \nu_3 R(Z_i^- X_j^+) R^\dagger + (\nu_4 + 1/2) Z_i X_j^-
\]

does not contain two-qubit interactions. The solution of this problem is non-trivial, mostly due to the many (nine in this case) parameters characterising a general unitary \(R\). To avoid searching solutions for such a system, we try a different approach to the problem. Let us denote with \(\tilde{H}\) a generator with the required properties: one that generates the same unitary as \(H_{\text{CNOT}}\) and contains only one-qubit interaction terms. Its general form will be

\[
\tilde{H} = h_0 + \sum_{\alpha=1}^3 (h_{1,\alpha}^0 \sigma_1^\alpha + h_{2,\alpha}^0 \sigma_2^\alpha).
\]  

(S4)

As discussed in section 2, for \(\tilde{H}\) to correctly generate the CNOT gate, it must commute with the principal generator \(H_{\text{CNOT}}\). Imposing this commutativity removes most of the parameters \(h_{1,\alpha}^0\), leaving us with the following simplified expression:

\[
\tilde{H} = h_0 + h_{1,3} Z_1 + h_{2,1} X_2.
\]  

(S5)

The only missing step is now to impose the eigenvalues of \(H_{\text{CNOT}} - \tilde{H}\) to be integer multiples of \(2\pi\). This gives the following system of equations:

\[
2\pi \nu_1 = (-\pi + 4h_0 - 4h_{1,3} - 4h_{2,1})/4,
2\pi \nu_2 = (+\pi + 4h_0 + 4h_{1,3} - 4h_{2,1})/4,
2\pi \nu_3 = (+\pi + 4h_0 - 4h_{1,3} + 4h_{2,1})/4,
2\pi \nu_4 = (-\pi + 4h_0 + 4h_{1,3} + 4h_{2,1})/4,
\]

(S6)

with \(\nu_i \in \mathbb{Z}\). The above system can be seen to have no solution for \(h_0, h_{1,3}, h_{2,1}\), therefore definitively proving that there is no rotation of the degenerate space, and integer parameters \(\nu_i\), that allow to generate the CNOT gate using only one-qubit interactions.

A3. Toffoli gate: derivation through conditions

We show here the details of how, using equation (1), we can obtain a family of Hamiltonian generators for the Toffoli gate, containing only single- and two-qubit interactions. This is a more detailed version of the analysis given in section 3.1.

The Toffoli gate can be written as

\[
U_{\text{toff}} = Z_1^+ + Z_1^- [Z_2^+ + Z_2^- (X_3^+ - X_3^-)],
\]  

(S7)

where we defined \(Z_i^\pm \equiv (1 \pm Z_i)/2\), and similarly for \(X_i^\pm\). The principal generator of \(U_{\text{toff}}\) is \(H_{\text{toff}} = \pi Z_1 Z_2 X_3\), that is, highlighting the three-qubit interaction term,

\[
H_{\text{toff}} = (1 \text{- and 2-qubit terms}) - \pi/8 Z_1 Z_2 X_3.
\]  

(S8)

We start by writing down the general parametrization of an Hamiltonian containing only one- and two-qubit interactions:
\[ \hat{H}_{\text{Toff}} = \hbar \mathbb{I} + \sum h_{ij} \sigma_i^a \sigma_j^b + \sum f_{ij}^{a\beta} \sigma_i^a \sigma_j^\beta. \]  

(S9)

Assuming for simplicity that \( \hat{H}_{\text{Toff}} \) does not contain \( Y \) components, and imposing the commutativity with \( \hat{H}_{\text{Toff}} \), we get the following expression

\[
\hat{H}_{\text{Toff}} = \hbar \mathbb{I} + h_{X3}^2 X_3 + h_{Z1}^2 Z_1 + h_{Z2}^2 Z_2 + f_{13}^{a2} Z_1 X_3 + f_{13}^{a2} Z_2 X_3 + f_{13}^{a2} Z_1 Z_2 \\
+ \left( f_{13}^{a2} X_1 + f_{13}^{a2} X_2 \right) (1 + X_3) + \left( f_{13}^{a2} X_2 + f_{13}^{a2} Z_3 \right) (1 + Z_1) \\
+ \left( f_{13}^{a2} X_1 + f_{13}^{a2} Z_2 \right) (1 + Z_2). \quad \text{(S10)}
\]

As can be directly verified, the above satisfies \([\hat{H}_{\text{Toff}}, \hat{H}_{\text{Toff}}] = 0\) while also containing only one- and two-qubit interactions, and no Pauli \( Y \) matrices. We could now directly try and find the set of parameters making the eigenvalues of \( \hat{H}_{\text{Toff}} = \hat{H}_{\text{Toff}} \) all integer multiples of \( 2\pi i \), but the associated calculations are made hard by the many parameters involved. However, it turns out that we can make further assumptions on the form of \( \hat{H}_{\text{Toff}} \) and still obtain a viable family of solutions. One of them, leading to a satisfying family of assumptions equation (S10) becomes

\[
\hat{H}_{\text{Toff}} = \hbar \mathbb{I} + h_{X3}^2 X_3 - \pi/8 (Z_1 + Z_2) (1 - X_3) + f_{13}^{a2} Z_1 Z_2 + f_{13}^{a2} (X_1 + X_2) (1 + X_3) + f_{13}^{a2} (Z_1 - Z_2) Z_3. 
\quad \text{(S11)}
\]

Finally, we impose that the generator is diagonal on the first two qubits, that is, \( f_{13}^{a2} = 0 \). Using this simplified expression, \( \hat{H}'_{\text{Toff}} = \hat{H}_{\text{Toff}} - \hat{H}_{\text{Toff}} \) becomes

\[
\hat{H}'_{\text{Toff}} = \pi/8 Z_1 Z_2 X_3 + (\hbar \mathbb{I} - \pi/8) 1 + (h_{X3}^2 + \pi/8) X_3 + (f_{13}^{a2} - \pi/8) Z_1 Z_2 + f_{13}^{a2} (Z_1 - Z_2) Z_3. \quad \text{(S12)}
\]

With the above simplified expression it is now possible to directly solve the eigenvalue problem. This results in the following family of solutions

\[
\hat{H}'_{\text{Toff}} = \frac{\pi}{8} \left[ 1 + 4 \left( \nu_1 + \nu_2 + 2 \nu_3 + \sqrt{(\nu_3 - \nu_4)^2} \right) - (Z_1 + Z_2) (1 - X_3) + X_3 (-2 - 8 \nu_1 + 8 \nu_2) \\
+ 4 Z_1 Z_2 \left( 1/4 + \nu_1 + \nu_2 - 2 \nu_3 - \sqrt{(\nu_3 - \nu_4)^2} \right) + (Z_1 - Z_2) Z_3 \sqrt{c(\nu_1, \nu_2, \nu_3, \nu_4)} \right], \quad \text{(S13)}
\]

with \( c(\nu_1, \nu_2, \nu_3, \nu_4) = (4 \nu_3 \nu_4)^2 - (1 + 4 \nu_3)^2 \), for all integer values of \( \nu_1 \) such that \( c(\nu_1, \nu_2, \nu_3, \nu_4) \geq 0 \). The corresponding spectrum of \( \hat{H}'_{\text{Toff}} = \hat{H}_{\text{Toff}} - \hat{H}_{\text{Toff}} \) is

\[
\lambda_1 = \lambda_2 = 2 \pi \nu_1, \\
\lambda_3 = \lambda_4 = 2 \pi \nu_2, \\
\lambda_5 = \lambda_6 = 2 \pi \nu_3, \\
\lambda_7 = \lambda_8 = 2 \pi (\nu_3 + |\nu_3 - \nu_4|),
\]

while the spectrum of \( \hat{H}'_{\text{Toff}} \) changes only in that \( \lambda_2 = 2 \pi (\nu_1 + 1/2) \). Consistently with this, \( \lambda_2 \) is also the eigenvalue corresponding to the non-degenerate eigenspace of \( \hat{H}_{\text{Toff}} \), while all the other eigenvalues correspond to eigenvectors orthogonal to this one. More specifically, we have

\[
|\lambda_1 \rangle = |0, 0, -\rangle, \quad |\lambda_2 \rangle = |1, 1, -\rangle, \\
|\lambda_3 \rangle = |1, 1, +\rangle, \quad |\lambda_4 \rangle = |0, 0, +\rangle, \\
|\lambda_5 \rangle = |1, 0\rangle \otimes N_{\alpha} \left[ (a - b)|0\rangle + |1\rangle \right], \\
|\lambda_6 \rangle = |0, 1\rangle \otimes N_{\alpha} \left[ (a + b)|0\rangle + |1\rangle \right], \\
|\lambda_7 \rangle = |1, 0\rangle \otimes N_{\alpha} \left[ (a + b)|0\rangle - |1\rangle \right], \\
|\lambda_8 \rangle = |0, 1\rangle \otimes N_{\alpha} \left[ (a - b)|0\rangle - |1\rangle \right],
\]

(S15)
understand how and why the derived full expressions for the projectors and the found solutions for the Toffoli gate. It is also interesting to note so that the sum of the projectors gives the identity as it should. On the other hand, multiplying by the appropriate three-qubit terms, we find

\[ P_1 + P_2 = \frac{1}{4}(1 + Z_1 Z_2)(1 - X_3), \quad P_3 + P_4 = \frac{1}{4}(1 - Z_1 Z_2)(1 + X_3). \]

It is easily verified from the above that

\[ P_1 + P_2 + P_3 + P_4 = \frac{1}{2}(1 + Z_2 Z_3), \quad P_5 + P_6 + P_7 + P_8 = \frac{1}{2}(1 - Z_2 Z_3), \]

so that the sum of the projectors gives the identity as it should. On the other hand, multiplying by the appropriate \( \nu_i \) factors, we get

\[ 2\pi [\nu_1(P_1 + P_2) + \nu_2(P_3 + P_4)] = (\ldots) + \frac{\pi}{2}(\nu_1 - \nu_2)Z_2X_3, \]

\[ 2\pi [\nu_3(P_5 + P_6) + \nu_4(P_7 + P_8)] = (\ldots) + \frac{\pi}{2}(\nu_3 - \nu_4)Z_2X_3 + \frac{\pi}{8}Z_2Z_3, \]

with the last identity holding for \( \nu_3 \neq \nu_4 \).

\[ a = \frac{|\nu_{34}|}{1 + \nu_{12}}, \quad b = \frac{\sqrt{|\nu_{34}|^2 - (\nu_{12} + 1)^2}}{1 + \nu_{12}}, \quad (S16) \]

It is worth noting that the orthogonality of these eigenvectors follows from the easily verified property of the above coefficients: \( a^2 - b^2 = 1 \). Furthermore, we note that \( c(\nu_1, \nu_2, \nu_3, \nu_4) \geq 0 \) cannot be satisfied unless \( \nu_3 \neq \nu_4 \). This in turn, looking at equation (S15), reveals that all the solutions are made possible by a non-trivial lifting of the degeneracy of the subspaces \( |0, 1\rangle \langle 0, 1| \) and \( |1, 0\rangle \langle 1, 0| \). Let us now try to understand how and why the derived \( H'_{\text{faff}} \) works. Let us use the notation \( P_i \equiv |\lambda_i\rangle \langle \lambda_i| \), and consider the projector over the last two eigenvectors. Highlighting the three-qubit terms, we find

\[ P_7 \simeq -N_{34}^2 \frac{Z_1 Z_2}{4} \left[ ((a + b)^2 - 1) \frac{Z_3}{2} - (a + b)X_3 \right], \]

\[ P_8 \simeq -N_{34}^2 \frac{Z_1 Z_2}{4} \left[ ((a - b)^2 - 1) \frac{Z_3}{2} - (a - b)X_3 \right]. \]

The term in the Hamiltonian to which these two projectors contribute is \( 2\pi \nu_{34}(P_7 + P_8) \), with \( \nu_{34} = \nu_3 + |\nu_3 - \nu_4| \). Recalling the definitions of \( a, b, N_3, N_6 \), we see that the coefficient of \( Z_1 Z_2 Z_3 \) vanishes, and the resulting expression becomes

\[ P_7 + P_8 \simeq Z_1 Z_2 X_3 \frac{1 + 4(\nu_1 - \nu_2)}{16|\nu_3 - \nu_4|}. \quad (S18) \]

Substitution of the appropriate values of \( \nu_i \) shows that the above term can be used to generate the three-qubit factor \( \pi/8 \ Z_1 Z_2 X_3 \), without introducing additional three-qubit factors. In table A1 are given the full expressions for the projectors and the found solutions for the Toffoli gate. It is also interesting to note that all of the above still holds if the \( X_i \) operators are replaced with \( Y_i \) operators. That is, the expressions found solving for the Toffoli, by simple substitution \( X_i \to Y_i \), also give a generator with only two-qubit interactions for the CCY gate (that is, the gate that applies \( Y \) to the third qubit conditionally to the first two qubits being in the \( |1\rangle \) state).
A different way to understand $\tilde{H}_{\text{Toff}}$ is to analyse the four two-dimensional subspaces on the main diagonal, exploiting the fact that $\tilde{H}_{\text{Toff}}$ acts diagonally on the first two qubits. Straightforward calculations lead to

$$\begin{align*}
(00)|\tilde{H}_{\text{Toff}}|00\rangle &= \pi \left[(\nu_1 + \nu_2) - (\nu_1 - \nu_2)X\right], \\
(01)|\tilde{H}_{\text{Toff}}|01\rangle &= 2\pi\nu_3 + \pi\nu_4(1 - \sigma_{00}), \\
(10)|\tilde{H}_{\text{Toff}}|10\rangle &= 2\pi\nu_3 + \pi\nu_4(1 - \sigma_{10}), \\
(11)|\tilde{H}_{\text{Toff}}|11\rangle &= \frac{\pi}{2} \left[(1 + 2(\nu_1 + \nu_2)) - (1 + 2\nu_2)X\right],
\end{align*}$$

where

$$
\sigma_{00} \equiv \frac{(1 + 4\nu_2)X + \sqrt{Z}}{4|\nu_4|}, \quad \sigma_{10} \equiv \frac{(1 + 4\nu_2)X - \sqrt{Z}}{4|\nu_4|}.
$$

It can be verified that for all values of $\nu_1, \nu_2, \nu_3, \nu_4$, the two-dimensional identity and $X$ are correctly generated in the $|00\rangle$ and $|11\rangle$ spaces, respectively. On the other hand, in the $|01\rangle$ and $|10\rangle$ spaces, the two-dimensional identity is generated as long as $\nu_3 \neq \nu_4$, as was also derived before. In particular, the class of solutions given by $\nu_1 = \nu_2 = \nu_3 = 0$ is

$$
\tilde{H} = \frac{\pi}{8} \left[ 1 + 4|\nu_4| - 2X_3 - Z_1 - Z_2 + (Z_1 + Z_2)X_3 + Z_1Z_2(1 - 4|\nu_4|) + (Z_2 - Z_1)Z_3 \sqrt{16\nu_4^2 - 1} \right],
$$

for all $\nu_4 \neq 0$. It is interesting to look at the explicit form of the exponentials generated by this class generators. Computing using equation (S20), we get

$$
\text{exp}(i\tilde{H}t) = \begin{pmatrix}
1 & 0 & 0 & 0 \\
0 & S(t, \nu_4) & 0 & 0 \\
0 & 0 & S(t, \nu_4) & 0 \\
0 & 0 & 0 & X(t)
\end{pmatrix},
$$

where

$$
S(t, \nu_4) = \begin{pmatrix}
a + b & c \\
c & a - b
\end{pmatrix},
$$

$$
a = \frac{1 + e^{2i\nu_4}}{2}, \quad \frac{1 - e^{2i\nu_4}}{8\nu_4},
$$

$$
b = \frac{(-1 + e^{2i\nu_4})\sqrt{16\nu_4^2 - 1}}{8\nu_4},
$$

and

$$
X(t) = \frac{1}{2} \begin{pmatrix}
1 + e^{2it} & (1 - e^{2it}) \\
(1 - e^{2it}) & (1 + e^{2it})
\end{pmatrix}.
$$

For large (in modulus) values of $\nu_4$, $a + b \to e^{2i\nu_4}$, $a - b \to 1$ and $c \to 0$, so that the exponential becomes

$$
\text{exp}(i\tilde{H}t) = \begin{pmatrix}
1 & e^{2i\nu_4} \\
e^{2i\nu_4} & 1 \\
1 & X(t)
\end{pmatrix},
$$

which very closely resembles the matrix obtained by exponentiating the principal generator $\mathcal{H}_{\text{Toff}} = \pi Z_1 Z_2 X_3$, that is

$$
\text{exp}(i\mathcal{H}_{\text{Toff}}t) = \begin{pmatrix}
1 & 0 \\
0 & 1 \\
0 & X(t)
\end{pmatrix}.
$$

A different solution derived from equation (S13) is

$$
\tilde{H}_{\text{Toff}} = \frac{9\pi}{8} + \frac{3\pi}{4}X_3 - \frac{\pi}{8}(Z_1 + Z_2) + \frac{\pi}{8}Z_1Z_2 + \frac{\pi}{8}(Z_1 + Z_2)X_3 - \frac{\sqrt{7}\pi}{8}(Z_1 - Z_2)Z_3.
$$
Moreover, it is worth noting that equation (S13) is only one possible family of solutions, and that different assumptions will lead to different ones. For example, a similar reasoning as above, starting however from the assumptions $I[f_{31}]^2 = f_{31}^2$ will lead to solutions such as [note the use of $(Z_1 + Z_2)$ terms here, making this solution not derivable from equation (S13)]

\[
\hat{H}_{\text{off}} = \frac{9\pi}{8} - \frac{7\pi}{8}X_3 + \frac{\sqrt{15}\pi}{8}Z_3 + \frac{\pi}{8}Z_1Z_2 \\
+ \frac{\pi}{8}(Z_1 + Z_2) \left( -1 + \frac{5}{2}X_1 + \frac{\sqrt{15}}{2}Z_1 \right). 
\]  

(A28)

**A4. Toffoli gate: an example of direct a posteriori derivation**

We will here show a line of thought that could have conceivably led to equation (S20) (in the case $\nu_4 = 1$), by direct analysis, and without using any of the tools shown in the paper. It will be useful to keep in mind the expressions of $Z_1 \pm Z_2$:

\[
Z_2 + Z_1 = \text{diag}(2, 2, 0, 0, 0, -2, -2), \\
Z_2 - Z_1 = \text{diag}(0, 0, -2, 2, 2, 0, 0).
\]

(S29)

Given that we want to generate a CC-X gate, and remembering that $\exp\left[\frac{\nu}{2}(1 - X)\right] = X$, it is reasonable to start building our Hamiltonian as

\[
\hat{H}_1 = -\pi \left( \frac{Z_1 + Z_2}{2} \right) \left( 1 - X_3 \right), 
\]

(S30)

which however will generate an X evolution both in the $|00\rangle$ and in the $|11\rangle$ sectors, while we want it only in the latter sector: $H_1 = \text{diag}(-X^-, 0, 0, X^-)$. We can remove the term in the $|00\rangle$ sector exploiting the sign difference introduced by $Z_1 + Z_2$, by directly adding an appropriate one-qubit interaction term:

\[
\hat{H}_2 = \frac{1}{2} \left[ \hat{H}_1 + \frac{\pi}{2} (1 - X_3) \right] \\
= \pi \text{ diag}(0, X^-/2, X^-/2, X^-), 
\]

(S31)

with $\exp(i\pi X^-) = X$. Equation (S31) now correctly reproduces the evolution on $|00\rangle$ and $|11\rangle$, but also wrongly evolves $|01\rangle$ and $|10\rangle$. To remove these additional terms while at the same time leaving the others unaffected we use the fact that $\exp(i(1 \pm \sigma)) = 1$, for any normalized vector of sigma matrices:

\[
\sigma \equiv \sum_{i=1}^{3} n_i \sigma_i \text{ with } n_1^2 + n_2^2 + n_2^2 = 1.
\]

To convert the central terms in equation (S31) into something like this we observe that we can rewrite the second term in the above equation as

\[
\pi/4(1 - X_3) = \pi/8(2 - 2X_3) = \pi/8(5 - 3 - 2X_3).
\]

(S32)

As $Z_1Z_2 = \text{diag}(1, -1, -1, 1)$, we substitute the above with $\pi/8(5 - 3Z_1Z_2 - 2X_3)$. This change affects only the central terms, converting the expression into: $\pi \text{ diag}(0, 1 - X/4, 1 - X/4, X^-)$. The reason this form is preferable is that we can now simply add a factor in the central terms to convert them into an expression of the form $1 - \sigma$. Adding an interaction of the form $\pi\alpha(Z_2 - Z_1)Z_3$ gives

\[
\pi \text{ diag}(0, 1 - X/4 - 2\alpha Z, 1 - X/4 + 2\alpha Z, X^-). 
\]

(S33)

For the central terms to exponentiate to the identity we need them to become of the form $1 - \sigma$ with normalized $\sigma$. This is easily achieved by choosing $\alpha = \pm \sqrt{15}/8$. The final expression is thus

\[
8/\pi \hat{H}_4 = - (Z_1 + Z_2)(1 - X_3) \pm \sqrt{15}(Z_2 - Z_1)Z_3 + (5 - 3Z_1Z_2 - 2X_3).
\]

(S34)

Note that instead of $\pi\alpha(Z_2 - Z_1)Z_3$ we could have equivalently chosen $\pi\alpha(Z_2 - Z_1)O_3$ for any $O_3 = aY_3 + bZ_3$ and $a^2 + b^2 = 1$. The above reasoning explains the origin of the weird $\sqrt{15}$ factor; it comes as the coordinate necessary to make the vector unitary: for $X/4 + xZ/4$ to be normalized, $x = \sqrt{15}$ must be satisfied.
Appendix B. Supervised learning approach

We here study more in depth the following problem: given a target gate $G$ and a parametrized Hamiltonian $H(\lambda) = \sum_k \lambda_k \sigma_k$, with $\lambda_k \in \mathbb{R}$ and $\sigma_k$ Hermitian operators, what is the set of parameters $\lambda_0$ such that $\exp(iH(\lambda_0)) = G$? We present a supervised learning approach to numerically solve this problem, considerably extending the ideas presented in reference [16]. Thanks to a number of numerical optimization techniques, and in particular the use of AD [20–23], we can explore a variety of different scenarios, optimizing over potentially hundreds of Hamiltonian parameters. On top of this, condition (1b) is used to further speed-up the numerical training, removing many interaction parameters that are known not to lead to the target gate.

B1. Supervised learning

Supervised learning is the task of inferring or approximating a function, given a set of pre-labeled data [18, 49]. A supervised learning algorithm starts with some model—a functional relation $g_\lambda$ parametrized by a set of parameters $\lambda$—and finds a $\lambda_0$ making $g_{\lambda_0}$ as close as possible to a target function $f$. To do this, a set of pre-labeled training data $\{(x_1,y_1),(x_2,y_2),\ldots\}$ is used, where here $y_2 = f(x_2)$ is the output that we want the algorithm to associate to the input $x_2$.

Among the most used supervised learning models are neural networks (NNs) [50, 51]. These are parametric non-linear models which play a prominent role in many machine learning tasks, such as dimensionality reduction, classification, and feature extraction [50, 51]. NNs have also recently proven useful for several problems in quantum many-body theory [52–59], quantum compilation [60], quantum stabilizer codes [61] and entanglement quantification [62].

A NN is trained by optimizing its parameters using a dataset of pre-labeled data. A common way to do this is use variations of a gradient-descent-based technique named SGD. Gradient descent algorithms aim to optimize a problem function $f(x)$, starting from an initial point $x_0$ and performing a number of small steps towards the direction of maximum slope (that is, $\nabla f(x)$). The optimal point $x_{opt}$ is thus obtained via a sequence of small perturbations of the point $x$, which starting from $x_0$ reaches the nearest local stationary point by following the slope. In the simplest version of the algorithm, the update rule is simply $x \rightarrow x - \eta \nabla f(x)$, with $\eta$ a small real parameter commonly referred to as learning rate. SGD, on the other hand, is suitable for a situation in which one is given a parametrized functional relationship of the form $f(x; w)$, and asked for a set of ‘parameters’ $w_0$ such that $f(x; w_0)$ is minimum (maximum) for all inputs $x$. Such a case can be handled via SGD, which in its simplest form involves picking a random $x_1$, executing a number of gradient descent iterations over $w$, then picking a new $x_2$ and iterating the procedure. The updating rule for SGD is therefore of the form

$$w \rightarrow w - \eta \nabla_w f(x;w).$$

(S35)
While standard gradient descent, being a local optimization algorithm, is liable to getting stuck in local minima, SGD can at least partially avoid this issue, in that generally a local minimum for an input $x$ is not a local minimum for a different input $x'$. Many variations of SGD are used in different circumstances. For example, in the so-called mini-batch SGD, instead of updating with a single input $x$, one uses a batch of inputs $\{x_1, \ldots, x_M\}$, and updates the parameters using the averaged gradient:

$$w \rightarrow w - \eta \frac{1}{M} \sum_{k=1}^{M} \nabla f(x_k; w) / M.$$  

More sophisticated updating rules are used to increase the training efficiency in different circumstances. Common techniques involve dynamically updating the learning rate, or using momentum gradient descent [37, 41] techniques.

To see how this class of optimization problems is relevant to us, consider the fidelity function $F$ defined as

$$F(\lambda, \psi) \equiv \langle \psi | G^\dagger \exp(iH(\lambda))G | \psi \rangle,$$  

with $G$ the target gate, $\lambda$ the set of parameters, and $\psi$ an input state. The gate design problem is then equivalent to finding $\lambda$ such that $F(\lambda, \psi)$ is maximised (that is, equal to 1) for all $\psi$. One possibility to solve this problem is to consider the average fidelity $F(\lambda)$, for which explicit formulas are known [63–65]. Standard optimization methods, like standard gradient descent or differential evolution, can be applied directly on $F(\lambda)$. This, however, reveals to be impractical, due to the complexity of the associated parameter landscapes. On the other hand, SGD allows to use a simple and efficient local maximisation.

**Figure B2.** Training histories for the Toffoli gate with only diagonal interactions. In each plot are reported the values of the nine network parameters during the training process, for each training epoch $t_e$. Each training process was left running until convergence to unit fidelity, therefore, the number of epochs in the horizontal axes differs for different training instances. The histories shown here correspond to training instances in which the parameters were initialised at various values, as seen from the leftmost values in each plot.
method, while at the same time being less prone to getting stuck in local maxima. This works particularly well in this case, because we know that the sets of parameters corresponding to the target gate are all and only those such that for all inputs $\psi$ the fidelity is unitary. Despite this, the supervised learning approach becomes more efficient when more states are used: using only states belonging to a basis would allow the optimizer to only operate on a small number of curves. Using an overcomplete set of states makes the algorithm much more robust and efficient.

A crucial step, efficiency-wise, in gradient descent algorithms, is the evaluation of the gradient. Numerically approximating the gradient, as done in previous works [16], is generally inefficient and scales badly with the number of optimized parameters. Here we will instead make use of the powerful technique of automatic differentiation (AD) [21, 23], described in appendix B2. AD dramatically improves the training efficiency, allowing to explore a richer variety of circumstances.

**B2. Backpropagation**

The gradient evaluation phase is efficiency-wise crucial for the training of a neural network. Computing the partial derivatives of the cost function with a standard method, like finite differences, has a complexity $O(N^3_w)$, with $N_w$ the number of parameters to differentiate [18]. This inefficiency can however be avoided.
using error backpropagation via AD. With this technique, the complexity of the gradient evaluation phase can be cut down to $O(N^2_\omega)$ [18]. This works by first decomposing the cost function of the model in terms of elementary operations, that is, functions the gradient of which is known analytically. In this way the computational graph representing the functional relation between input and output is built. A computational graph is a directed acyclic graph, whose nodes represent the operations, and edges the flowing direction of inputs into outputs (see figure B1). Once the computational graph is built, the gradients with respect to the model parameters can be computed efficiently. This happens in two stages, as schematically illustrated in figure B1. At first, every node of the computational graph is progressively computed, starting from the inputs (the current values of the model parameters) up to the final value of the error function. During this process, the intermediate values of the elementary operations are cached. This is the so-called feed-forward phase. The second phase (so-called backpropagation phase) starts from the output, and consists of progressively computing the gradients of the error function with respect to the independent variables.

To better understand AD, let us consider a simple example. Suppose the error function of the model is of the form $g(w) \equiv f(f^{(2)}(f^{(1)}(w)))$, where $w$ is a set of parameters, and $f^{(i)}$ are intermediate 'elementary' functions, the gradients of which are supposed to be known analytically. Making use of the chain rule, the gradient of $g$ reads

$$\nabla g(w) = \sum_k \partial_k f(y^{(2)}) \nabla f^{(2)}_k(y^{(1)}),$$

(S37)
Figure B5. Fidelity $F_\lambda (\psi)$ vs variations of $\lambda$, for different test states, for the Toffoli gate. The five test states $\psi$ are sampled randomly. (a) Global relative variations of $\lambda$, that is, plotting the fidelity against $\alpha \lambda$ for $0.9 \leq \alpha \leq 1.1$. Note that this is equivalent to studying how the fidelity changes with respect to uncertainties in the evolution time, that is, how much does $\exp(iHt')$ differ from $\exp(iHt)$. (b) Same as (a) but with $0 \leq \alpha \leq 1.2$. (c) Plot of $F_\lambda (\psi)$ against absolute variations of a single element of $\lambda$, in this case the first one, i.e. we take $\lambda_1 \in [-10,10]$. (d) Like (c) but for $\lambda_2$.

Figure B6. Fidelity $F_\lambda (\psi)$ vs variations of $\lambda$, for different test states, for the Fredkin gate. The five test states $\psi$ are sampled randomly. (a) Global relative variations of $\lambda$, that is, plotting the fidelity against $\alpha \lambda$ for $0.9 \leq \alpha \leq 1.1$. Note that this is equivalent to studying how the fidelity changes with respect to uncertainties in the evolution time, that is, how much does $\exp(iHt')$ differ from $\exp(iHt)$. (b) Same as (a) but with $0 \leq \alpha \leq 1.2$. (c) Plot of $F_\lambda (\psi)$ against absolute variations of a single element of $\lambda$, in this case the first one, i.e. we take $\lambda_1 \in [-10,10]$. (d) Like (c) but for $\lambda_2$.

where $y^{(2)} = f^{(2)}(f^{(1)}(w))$ and $y^{(1)} = f^{(1)}(w)$. During the feed-forward phase the values of $y^{(1)}$ and then $y^{(2)}$ are progressively computed and cached. Using $y^{(2)}$, and the known expression for $\partial_j f_j$, $\partial_j f_j(y^{(2)})$ is then efficiently computed. The process continues by evaluating $\nabla f_k^{(2)}$, which is written as

$$\nabla f_k^{(2)}(y^{(1)}) = \sum_j \partial_j f_k^{(2)}(y^{(1)}) \nabla f_j^{(1)}(w).$$

(S38)

Again, being $y^{(1)}$ already computed during the feed-forward, $\partial_j f_j^{(2)}(y^{(1)})$ is readily computed. The last component needed for the full gradient is $\partial_j f_j^{(1)}(w)$, all parts of which are known. This method therefore
allows to efficiently evaluate numerical the gradient of complicated functions, without approximating the derivatives.

In the context of training neural networks, the function to be derived is the cost function of the network, that is, roughly speaking, the (Euclidean) distance between the result obtained for an input and the corresponding training output. For the gate design problem, we will use another notion of distance between output obtained and output expected. For quantum states, the fidelity between these turns out to work well.

B3. Implementation details

We used Python as language of choice for the implementation of the supervised learning. Being Python language of widespread use in the machine learning community, many libraries and frameworks are available to build computational graphs over which AD can be used. In particular, we used Theano [38], together with the QuTiP library for the simulation of the dynamics of quantum systems [66, 67].

Our implementation allows the training of an arbitrary target gate, parametrized via a time-independent Hamiltonian $\mathcal{H}(\lambda)$. The parametrization is completely arbitrary (provided the dependence on the parameters is linear), so that the Hamiltonian can be chosen as $\mathcal{H}(\lambda) = \sum_i \lambda_i A_i$ for any set of matrices $A_i$ and number of parameters $\lambda_i$. This is made possible by the flexibility of AD, which allows to automatically build an efficiently differentiable computational graph, without needing to hardcode the structure of the Hamiltonian.

The goal of the algorithm is, given a target gate $G$ and a parametrization for the Hamiltonian $\mathcal{H}(\lambda)$, find the $\lambda_0$ such that $\exp(i\mathcal{H}(\lambda_0)) = G$. We use for the purpose mini-batch SGD with momentum. The mini-batch version of SGD involves computing the gradient, at every iteration, averaging over the gradients computed for a number of states. Making such batches of states larger or smaller allows to enhance or decrease the variance of the gradients with respect to the input state. The use of momentum [37, 41] involves using a modified version of equation (S1). The updating rule is instead given by

$$
\nu \rightarrow \gamma \nu + \eta \nabla_{\lambda} F(\psi, \lambda), \\
\lambda \rightarrow \lambda + \nu,
$$

where here $\eta$ is the learning rate and $\gamma$ the momentum. The use of the auxiliary parameter $\nu$ during the training discourages sudden changes of direction, and can make the training significantly more efficient [41].

![Figure B7.](image-url)
Figure B8. Training results for the Fredkin gate with a generator containing one-qubit interactions and two-qubit interactions of the form $J_{ij}(X_iX_j + Y_iY_j)$. Every point shows the final fidelity obtained at the end of a training procedure. The hyperparameters, as well as the total number of training iterations, are kept the same in all the training instances shown here. The initial parameters’ values are the same within each horizontal sector, but changed between different sectors. The initial parameters’ values within each sector have been chosen as all equal to $c$ (that is, $\lambda_i = c$ for all $i$). The values of $c$ are 0, 1, ..., 10, with $c = 0$ in the leftmost sector and $c = 10$ in the last to rightmost one. The rightmost sector contains the results of training attempts with the initial values chosen at random (that is, with $\lambda_i$ sampled according to the uniform normal distribution, independently for each $i$). The greatest reached fidelity, obtained with initial parameters’ values $\lambda_i = 3$, is $F \approx 0.94$.

While the cost function $F$ is always real, some of the intermediate calculations needed to compute it involve complex numbers. While this poses no fundamental problems, many of the machine learning (ML) libraries do not support AD over functions with complex inputs or outputs. We worked around this problem using a similar trick to the one reported in [68]. In particular, to use the existing framework, we mapped the problem into one involving only real numbers. To do this, we map complex matrices into real ones via the bijection $A \mapsto \operatorname{Re}(A) \equiv 1 \otimes A_R - i\sigma_y \otimes A_I$, where $A_R$ and $A_I$ are the real and imaginary parts of $A$, respectively. At the same time, state vectors are to be mapped to $\Psi \mapsto \operatorname{Re}(\Psi) \equiv (\Psi_R, \Psi_I)^T$. It is easy to verify that with this mapping $A\Psi \mapsto \operatorname{Re}(A\Psi) = \operatorname{Re}(A)\operatorname{Re}(\Psi)$, so that all calculations can be equivalently be carried out with the real versions of matrices and vectors.

More specifically, the employed algorithm involves the following steps:

(a) Choose an initial set of parameters $\lambda$ (randomly, or specific values if one has an idea of where a solution might be). A number of other hyperparameters have to be decided at this step, depending on the exact SGD method used. In particular, for mini-batch SGD with momentum and decreasing learning rate, one has to decide the momentum $\gamma$, the initial value of $\eta$, the rate at which $\eta$ decreases during the training, and the size $N_b$ of the batches of states used for every gradient descent step.

(b) Repeat the following loop $N_e$ times, or until a satisfactory result is obtained. Each such iteration is conventionally named an epoch. Another hyperparameter to be chosen beforehand is the number of training states $N_t$ to be used in each epoch. Once this is fixed, every epoch will involve a number $N_b/N_t$ of gradient descent steps, each one using $N_e$ states for a single gradient calculation. $N_e$ random training states are sampled, to be used during the epoch.

1. Pick $N_b$ of the $N_e$ training states.
2. Forward-propagate each state of the sample, and then backpropagate the gradients, thus computing the average gradient over the mini-batch $\nabla_{\lambda} F(\lambda)$.
3. Update the coupling strengths $\lambda$ as per equation (S5).
4. Return to point (1).

B4. Results
A sample of training results for Toffoli, Fredkin, and ‘double Fredkin’ gates are given in figure 1. In figures B2–B4 are shown the training histories of the parameters for eight different solutions for Toffoli, Fredkin and double Fredkin, respectively. These illustrate how quickly the networks converge for different
initial values of the parameters. In all the shown cases the target gates are obtained with unit fidelity up to numerical precision (that is, all fidelities are between $1 \times 10^{-16}$ and 1). Different sets of optimization hyperparameters are found to give acceptable solutions. For the trainings shown in this paper we used a dynamically updated learning rate given, for the \( k \)th epoch, by $\eta = 1/(1 + k \alpha)$ with the decay rate $\alpha = 0.005$. The other hyperparameters were chosen as $\gamma = 0.5$, $N_b = 2$, $N_{tr} = 200$. Different initial values for the parameters were tested, but in most cases we started the training with either vanishing or random (following a normal distribution) parameters. For the training of the four-qubit gate we found the network to converge sooner to a solution when the parameters were initialised to a positive value (often with all parameters initialised to 4).

In figures B5–B7 we report the behaviour of the fidelity upon changes of the learnt Hamiltonian parameters, for Toffoli, Fredkin and double Fredkin gates, respectively. As shown in these plots, the stability of the implemented gates with respect to variations of time and interactions values greatly varies between different solutions, as well as between different parameters in the same solutions.
Training results for the Toffoli gate with a generator containing all one-qubit interactions and two-qubit interactions of the form $J^{(1)}_{ij} X_i X_j + J^{(2)}_{ij} Y_i Y_j$. The initial conditions are chosen as in figure B8. The maximum fidelities obtained are $F \simeq 0.98$, obtained in the $c = 4$ and $c = 5$ sectors.

To assess the feasibility of non-trivial gates in more restrictive experimental scenarios, we performed a systematic analysis of the reachability of Fredkin and Toffoli gates when allowing only for single-qubit and $X_i X_j + Y_i Y_j$ two-qubit interactions, and in the less restrictive setting of allowing for all $X_i X_j$ and $Y_i Y_j$ interactions. The results are shown in figures B8–B11. For the Fredkin gate, in the more restrictive $XX$ interactions setting, the biggest fidelity obtained was $F \simeq 0.94$, while when allowing for all $XX$ and $YY$ interactions the maximum fidelity obtained was $F \simeq 0.999$. For the Toffoli gate, the maximum fidelity obtained in the $XX$ scenario was $F \simeq 0.94$ as well, while when allowing for all $XX$ and $YY$ interactions the best training results corresponded to $F \simeq 0.98$. To have more consistent results, in all the training instances shown here all the hyperparameters, except for the interaction parameters’ initial values, were chosen to have the same value. In particular, each training instance was run for 200 epochs, each one using 200 random quantum states as inputs, divided in batches of two elements. This choice of hyperparameters is mostly empirical, and it is possible for different values to provide better results.

The above provides further evidence for the flexibility of the supervised learning approach, which can produce solutions with good fidelities even in more restrictive scenarios, closer to the capabilities of state of the art experimental architectures. Furthermore, the values of the interaction strengths for many of the presented solutions are found to be compatible with the capabilities of state of the art circuit-QED architectures with gate times of the order of tens of nanoseconds [26, 27]. For instance, the Hamiltonian reported in the supplementary information (https://stacks.iop.org/NJP/22/065001/mmedia) of reference [26], after a suitable adiabatic elimination of the field used to mediate the interaction among various superconducting qubits, returns a model that is very close to the ansatz used in equation (10).

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